

# An introduction to Construction Schemes

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$F \subseteq [\omega_1]^{<\omega}$  is a construction scheme of type

$\{(m_{k+1}, n_{k+1}, r_{k+1})\}_{k \in \omega}$  if there is a decomposition

$F = \bigcup_{k \in \omega} F_k$  such that for all  $k \in \omega$ :

1)  $F_0 = [\omega_1]'$

2)  $F$  is cofinal in  $[\omega_1]^{<\omega}$

3) If  $F \in \mathcal{F}_k$ , then  $|F| = m_k$

4) If  $F, E \in \mathcal{F}_k$ , then  $F \cap E \subseteq F, E$

5) For all  $F \in \mathcal{F}_{k+1}$ , there are  $\{F_i | i < n_{k+1}\} \subseteq \mathcal{F}_k$  and  $RC(F)$  such that:

a)  $F = \bigcup_{i < n_{k+1}} F_i$

b)  $\{F_i | i < n_{k+1}\}$  is a  $\Delta$ -system with root  $RC(F)$

c)  $|RC(F)| = r_{k+1}$

d)

$$RC(F) \subset F_0 \setminus RC(F) \subset F_1 \setminus RC(F) \subset \dots \subset F_{n_{k+1}-1} \setminus RC(F)$$

Each  $F \in \mathcal{F}_{k+1}$  looks like this:

$RCF$ )

$F_0 \setminus RCF$ )

$F_1 \setminus RCF$ )

...



Size  $r_{k+1}$

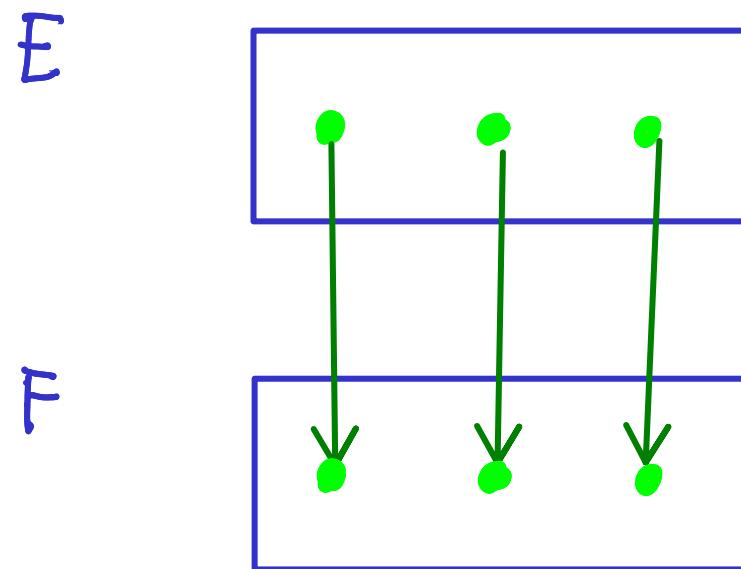
$n_{k+1}$  pieces

Theorem (Todorcevic)

Construction Schemes exist (for  
any type)

Def

Let  $E, F \in \omega_1^{\text{reg}} <^\omega$  of the same size. Denote by  $\varphi_{EF}$  the only increasing bijection from  $E$  to  $F$ .



Applications of construction schemes  
are often obtained as follows:

Assume we want to build certain  
kind of structure of size  $w$ ,

Fix  $F$  a construction scheme

1) Find  $\mathbb{P}$  a forcing consisting of finite approximations of the desired structure

2) For every  $F \in \mathcal{F}$ , find some  $p_F \in P$   
such that:

a) If  $E \in \mathcal{F}$  and  $E \subseteq F \Rightarrow p_F \leq p_E$

( $p_F$  is a better approximation than  $p_E$ )

b) If  $E, F \in \mathcal{F}_k$ , then  $p_F$  and  $p_E$  are  
computable. Moreover, they must be  
"coherent".

3) Look at the structure generated  
by {P<sub>F</sub>|FεF}

4) ... hope for the best!

We will construct a gap and  
an Aronszajn tree with construction  
schemes. I learn this constructions  
In the thesis of Fulgencio López.

Constructing a gap

Let  $A, B \subseteq \omega$

We say that  $A \subseteq^* B$  ( $A$  is almost contained in  $B$ ) if  $A \setminus B$  is finite

1) We say that  $\mathcal{L} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is a pregap if for every  $\alpha < \delta$ :

$$A_\alpha \subseteq^* A_\beta \subseteq^* B_\beta \subseteq^* B_\alpha$$

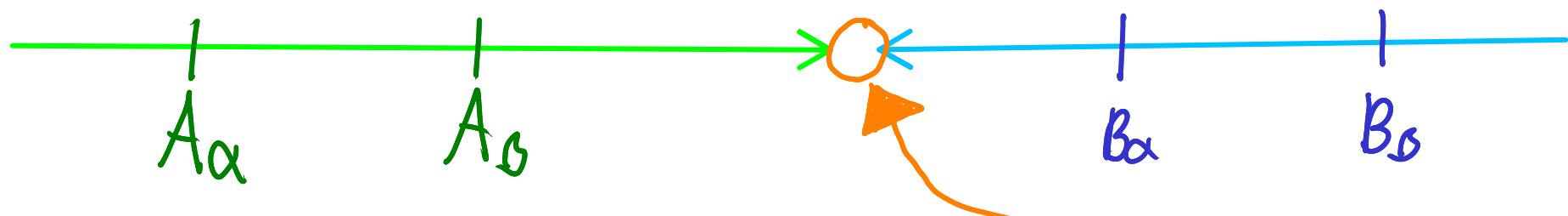


2) A pregap  $\mathcal{G} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is a gap if there is no  $C$  such that:

$$A_\alpha \subseteq^* C \subseteq^* B_\alpha \text{ for all } \alpha < \omega_1$$

2) A pregap  $\mathcal{G} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is a gap if there is no  $C$  such that:

$$A_\alpha \subseteq^* C \subseteq^* B_\alpha \text{ for all } \alpha < \omega_1$$



Nothing can be placed here

Gaps are very concrete examples  
that  $P(\omega)/fin$  is not complete

Theorem (Hausdorff)

There is a gap

# Theorem (Hausdorff)

There is a gap

The usual proof is by a recursion of size  $\omega_1$ . We will prove the Theorem using construction schemes

Knowledge of gaps is fundamental  
when we are trying to embed  
structures into  $P(\omega)/fin$

Gaps represent possible obstructions  
we may encounter.

For example, assume we have a linear order  $L = \{p_\alpha \mid \alpha \in \omega_2\}$  and we want to construct an embedding  $F: L \rightarrow P(\omega)/fin$  (of course, we are in a world where CH fails).

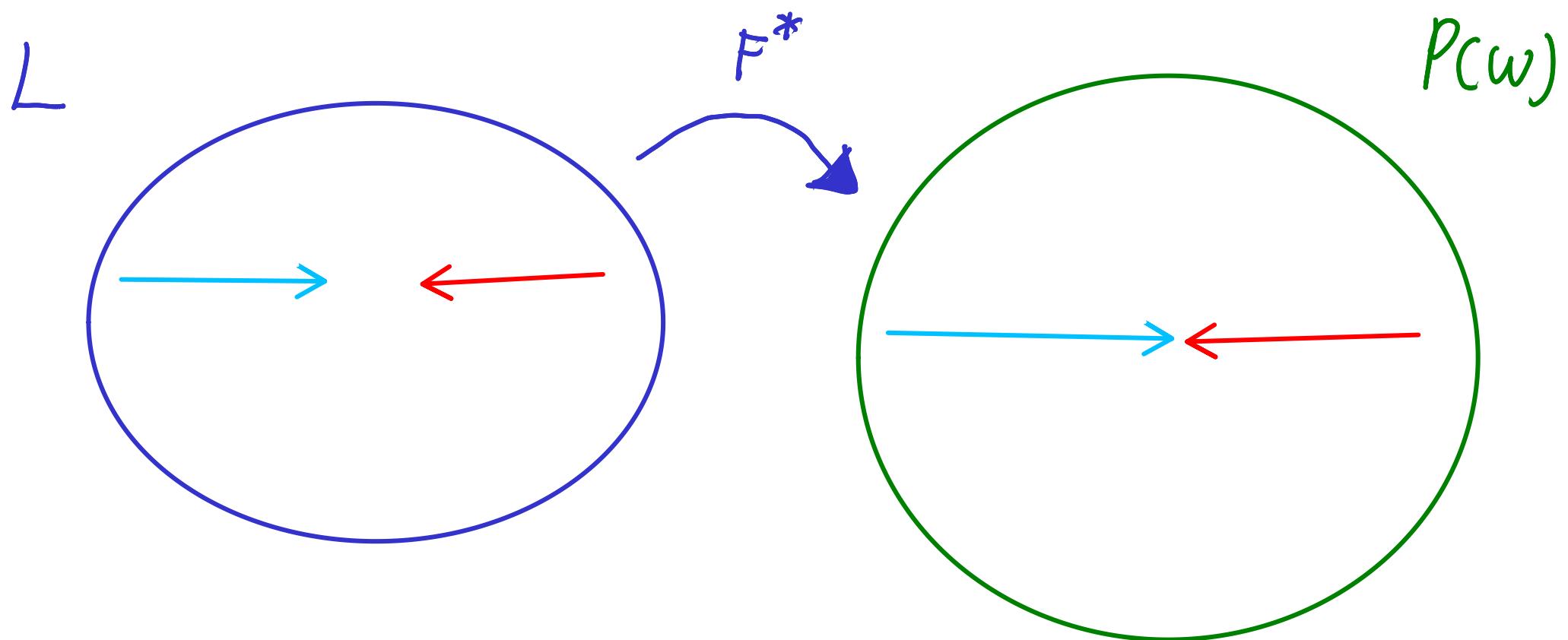
We try to define  $F$  recursively.

It could happen that in the  $\omega_1$  step, we have a partial embedding

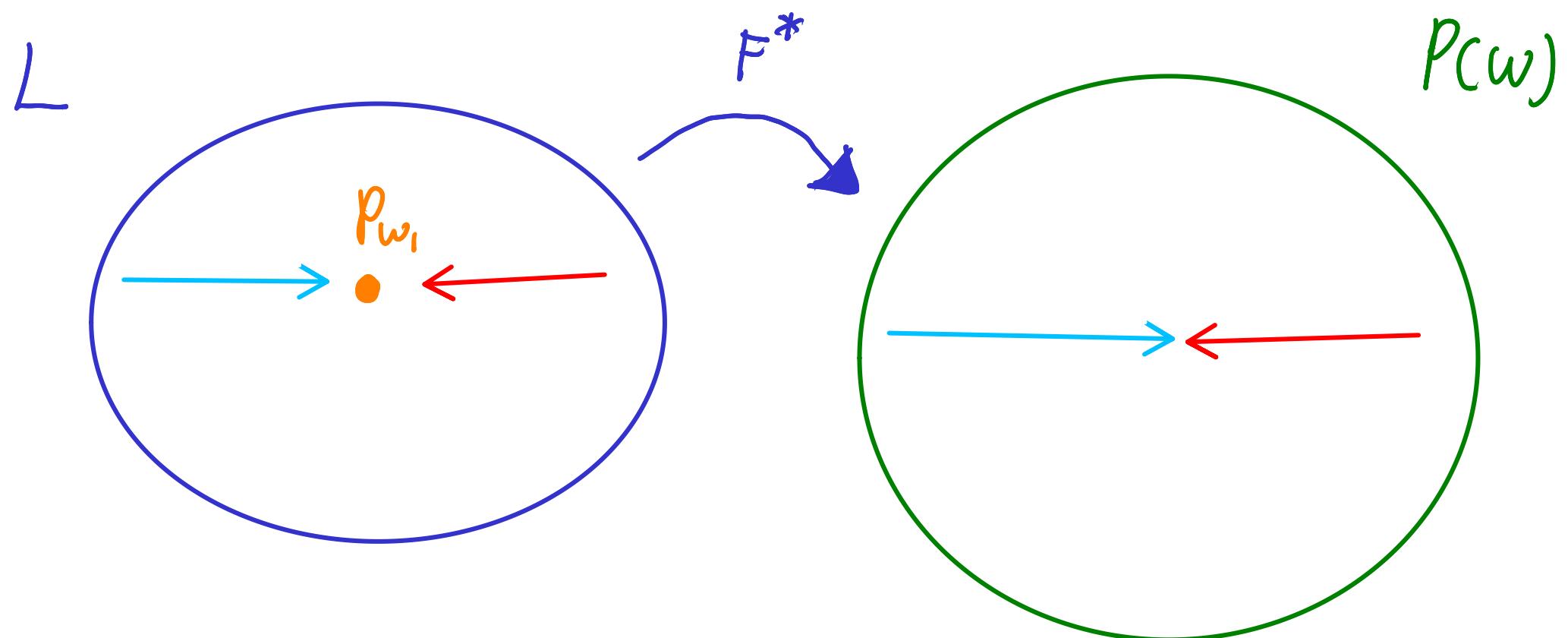
$$F^* : \{p_\alpha \mid \alpha < \omega_1\} \longrightarrow P(\omega)$$

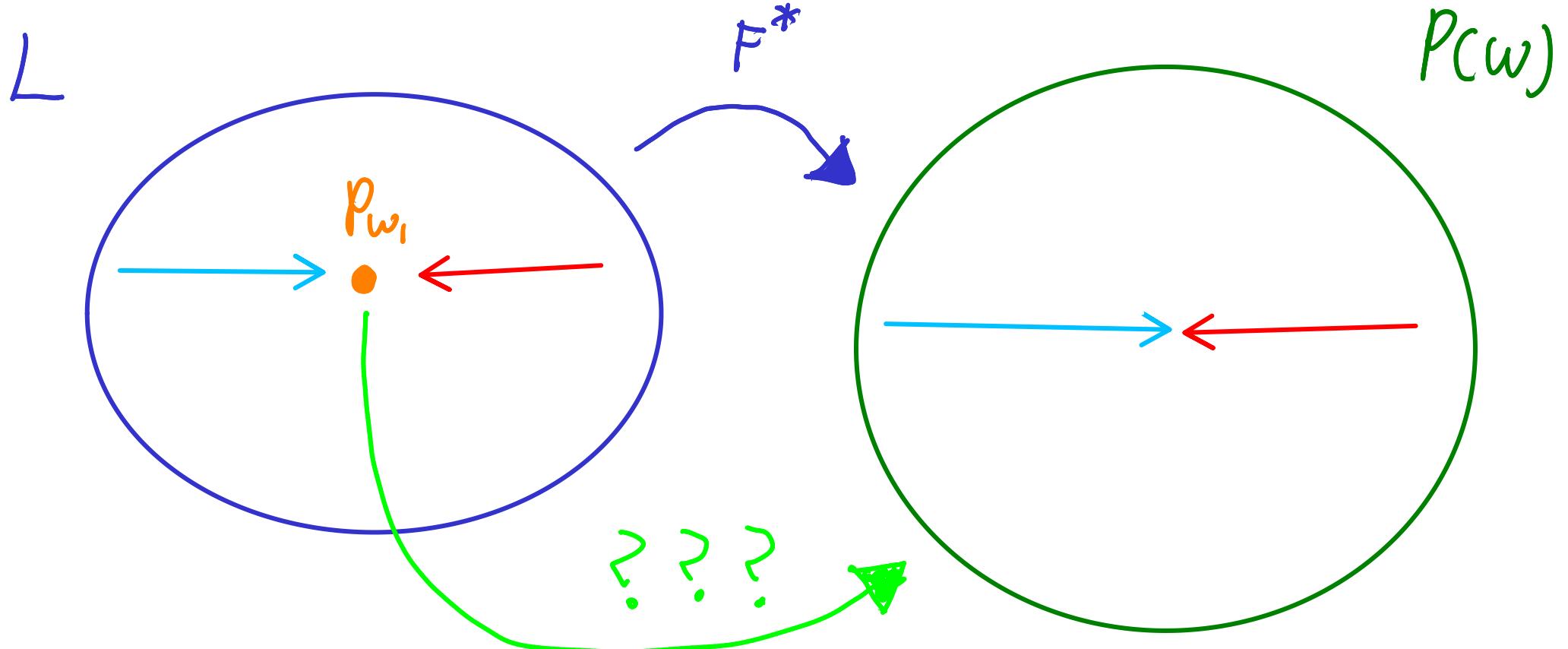
such that:

1) There are  $X, Y \in [w, j^\omega]$  such that  
 $(F^*[P_{\lambda \in X}], F^*[P_{\lambda \in Y}])$  is  
a gap



2)  $p_\alpha < p_{w_1} < p_\beta$  for all  $\alpha \in X$  and  $\beta \in Y$





In here,  $F^*$  can not be extended  
to an embedding, so our attempt  
to embed  $L$  into  $P(\omega)/fin$  failed.

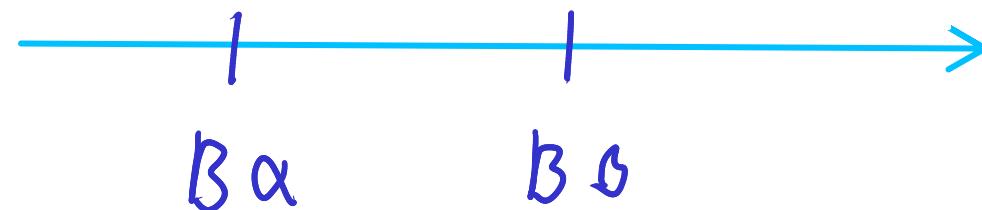
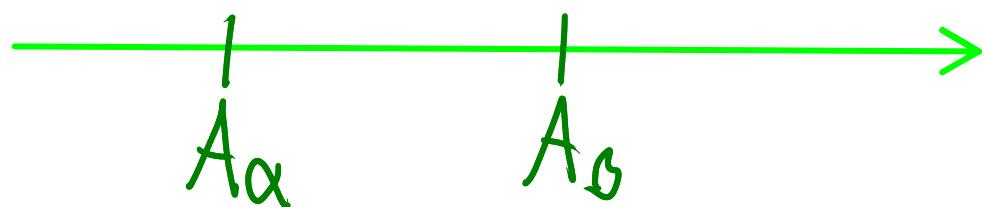
It will be convenient to reformulate  
the notion of gap

1) We say that  $\mathcal{G} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is  
a pregap if for every  $\alpha < \delta$ :

$$1) A_\alpha \subseteq^* A_\beta$$

$$2) B_\alpha \subseteq^* B_\beta$$

$$3) A_\alpha \cap B_\alpha = \emptyset$$



2) A pregap  $\mathcal{G} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is a gap if there is no  $C$  such that:

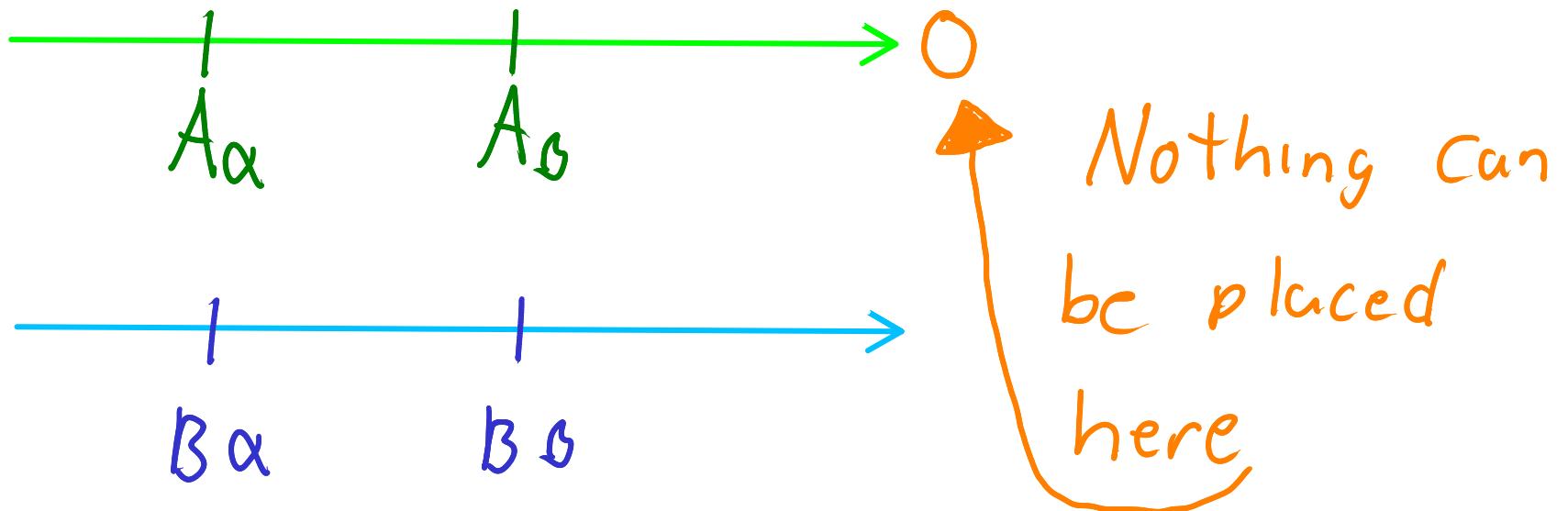
1)  $A_\alpha \subseteq^* C$  for all  $\alpha < \omega_1$

2)  $B_\alpha \cap C$  is finite for all  $\alpha < \omega_1$

2) A pregap  $\mathcal{E} = (\langle A_\alpha \rangle_{\alpha < \omega_1}, \langle B_\alpha \rangle_{\alpha < \omega_1})$  is a gap if there is no  $C$  such that:

1)  $A_\alpha \subseteq^* C$  for all  $\alpha < \omega_1$ ,

2)  $B_\alpha \cap C$  is finite for all  $\alpha < \omega_1$ .



Both definitions of gaps are essentially the same.

From now on, every mention of "gap" refers to the second meaning

Fix  $\mathcal{F}$  a  $\langle (m_{k+1}, 2, r_{k+1}) \rangle_{k \in \omega}$  construction scheme

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This means we are always gluing two things of  $\mathcal{F}_k$  to produce one in  $\mathcal{F}_{k+1}$



Let  $\mathbb{P}$  be the set of all  $p = (X, n, \Delta_p, \mathcal{B}_p)$ :

1)  $X \in [w, j^{<\omega}]$  and  $n \in \omega$

$X$  will be referred as the domain of  $p$  ( $X = \text{dom}(p)$ ) and  $n$  as the height of  $p$  ( $n = \text{ht}(p)$ )

Let  $\mathbb{P}$  be the set of all  $p = (X, n, A_p, B_p)$ :

1)  $X \in [\omega_1]^{<\omega}$  and  $n \in \omega$

2)  $A_p = \langle A_\alpha^p \rangle_{\alpha \in X}, B_p = \langle B_\alpha^p \rangle_{\alpha \in X}$

Let  $\mathbb{P}$  be the set of all  $p = (X, n, A_p, B_p)$ :

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2)  $A_p = \langle A_\alpha^p \rangle_{\alpha \in X}, B_p = \langle B_\alpha^p \rangle_{\alpha \in X}$

3)  $A_\alpha^p, B_\alpha^p \subseteq n$  for  $\alpha \in X$

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4)  $A_\alpha^p \cap B_\alpha^p = \emptyset$  for  $\alpha \in X$

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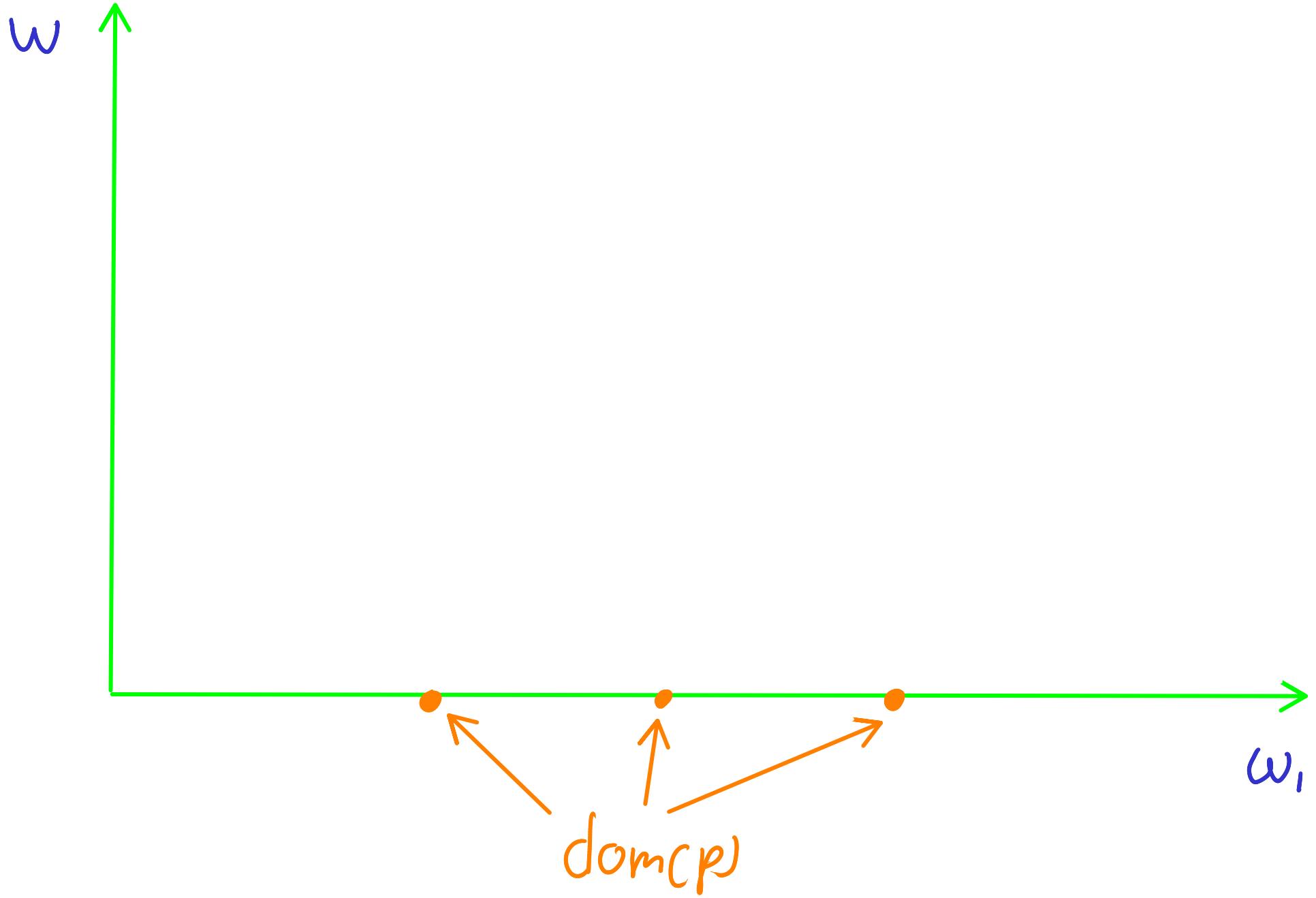
1)  $X \in [\omega_1]^{\leq \omega}$  and  $n \in \omega$

2)  $A_p = \langle A_\alpha^p \rangle_{\alpha \in X}, B_p = \langle B_\alpha^p \rangle_{\alpha \in X}$

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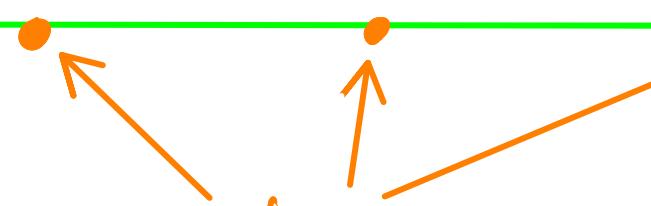


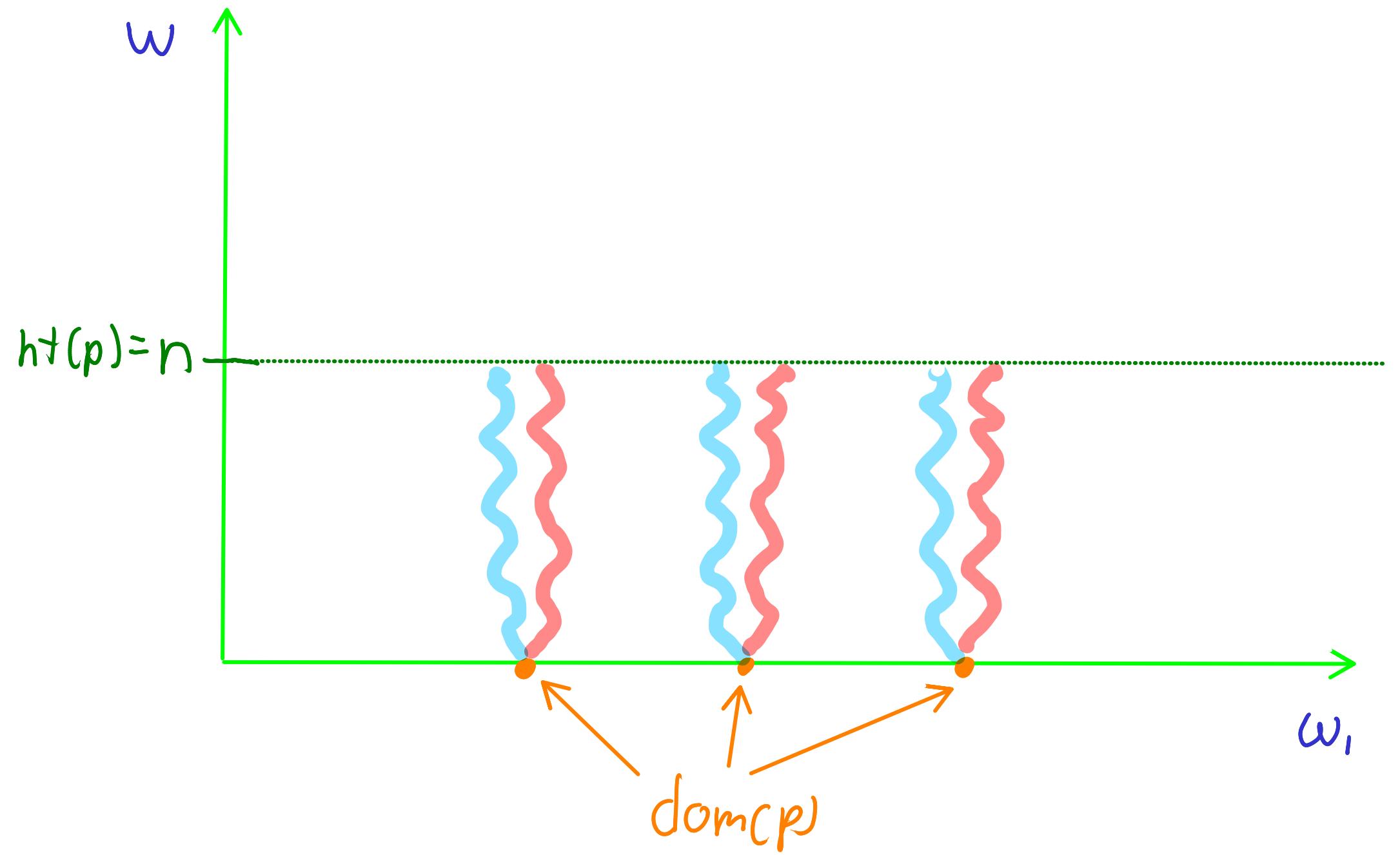
$w$

$ht(p) = n$

$\omega_1$

$dom(p)$





Let  $p = (X_p, n_p, \lambda_p, \mathcal{B}_p)$  and  $q = (X_q, n_q, \lambda_q, \mathcal{B}_q)$

Define  $p \leq q$  ( $p$  is a better approximation than  $q$ ) if:

Let  $p = (X_p, n_p, A_p, \mathcal{B}_p)$  and  $q = (X_q, n_q, A_q, \mathcal{B}_q)$

Define  $p \leq q$ :

1)  $X_q \subseteq X_p$  and  $n_q \leq n_p$

2) If  $\alpha \in X_q$ , then:

$$a) A_\alpha^q = A_\alpha^p \cap n_q$$

$$b) B_\alpha^q = B_\alpha^p \cap n_q$$

3) If  $\alpha, \beta \in X_q$  with  $\alpha < \beta$ , then:

a)  $A_\alpha^P \setminus n_q \subseteq A_\beta^P$

b)  $B_\alpha^P \setminus n_q \subseteq B_\beta^P$

Define  $p \leq q$ :

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3) If  $\alpha, \beta \in X_q$  with  $\alpha < \beta$ , then:

$$A_\alpha^p \setminus n_q \subseteq A_\beta^p \quad \text{and} \quad B_\alpha^p \setminus n_q \subseteq B_\beta^p$$

We will now define  $\langle N_k | k \in \omega \rangle$  and  $\langle p_F | F \in \mathcal{F} \rangle$   
such that:

I)  $\langle N_k \rangle_{k \in \omega}$  is an increasing sequence  
of natural numbers

2) If  $F \in \mathcal{F}_k$ , then:

a)  $\text{dom}(p_F) \supseteq F$

b)  $h^t(p_F) = N_k$

2) If  $F \in \mathcal{F}_k$ , then:

a)  $\text{dom}(p_F) \supseteq F$

b)  $ht(p_F) = N_k$

For convenience, define:

$$A_\alpha^F = A_\alpha^{p_F} \quad \text{and} \quad B_\alpha^F = B_\alpha^{p_F}$$

For  $\alpha \in F$

3) If  $F, E \in \mathcal{F}$  and  $E \subseteq F$ ,

then  $p_F \leq p_E$

(If  $F$  is bigger than  $E$ , then

$p_F$  is a better approximation

than  $p_E$ )

4) Let  $E, F \in \mathcal{F}_{k_1}$

a)  $p_E$  and  $p_F$  are compatible

(there is  $r \in P$  with  $r \leq p_E, p_F$ )

4) Let  $E, F \in \mathcal{F}_k$ ,

- a)  $P_E$  and  $P_F$  are compatible
- b) If  $\alpha \in F$  and  $\varphi = \varphi_{FE}$ , then:

$$A_\alpha^F = A_{\varphi(\alpha)}^E$$

$$B_\alpha^F = B_{\varphi(\alpha)}^E$$

1)  $\langle N_k \rangle_{k \in \omega}$  is an increasing sequence

2) If  $F \in \mathcal{F}_k$ , then:  $\text{dom}(p_F) \supseteq F$  and  
 $h(p_F) = N_k$

3) If  $F, E \in \mathcal{F}$  and  $E \subseteq F$ , then  $p_F \leq p_E$

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$$A_\alpha^F = A_{\varphi(\alpha)}^E \quad \text{and} \quad B_\alpha^F = B_{\varphi(\alpha)}^E$$

We define this items recursively

base case:  $k=0$

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Let  $N_0 = 2$ .

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Let  $N_0 = 2$ .

Choose  $F \in \mathcal{F}_1$ , say  $F = \{\alpha\}$ . Define:

$p_F = (\{\alpha\}, 2, \lceil A_\alpha^F \rceil, \lceil B_\alpha^F \rceil)$  where:

base case:  $k=0$

Let  $N_0 = 2$ .

Choose  $F \in \mathcal{F}_1$ , say  $F = \{\alpha\}$ . Define:

$p_F = (\{\alpha\}, 2, \lceil A_\alpha^F \rceil, \lceil B_\alpha^F \rceil)$  where:

$$A_\alpha^F = \{0\} \quad \text{and} \quad B_\alpha^F = \{1\}$$

recursive step:

Assume we defined  $N_k$  and  $p_E$  for  
 $E \in \mathcal{F}_k$ .

We will now define  $N_{k+1}$  and  $p_F$  for  
 $F \in \mathcal{F}_{k+1}$

Let  $N_{k+1} = N_k + 5$ .

Pick  $F \in \mathcal{F}_{k+1}$ . We know there are  
 $F_0, F_1 \in \mathcal{F}_k$  and  $R(F)$  such that:

$$1) F = F_0 \cup F_1$$

$$2) F_0 \cap F_1 = R(F)$$

$$3) |R(F)| = r_{k+1}$$

$$4) R(F) \subset F_0 \setminus R(F) \subset F_1 \setminus R(F)$$

Pick  $F \in \mathcal{F}_{k+1}$ . We know there are  
 $F_0, F_1 \in \mathcal{F}_k$  and  $RC(F)$  such that:

$F$



$RC(F)$

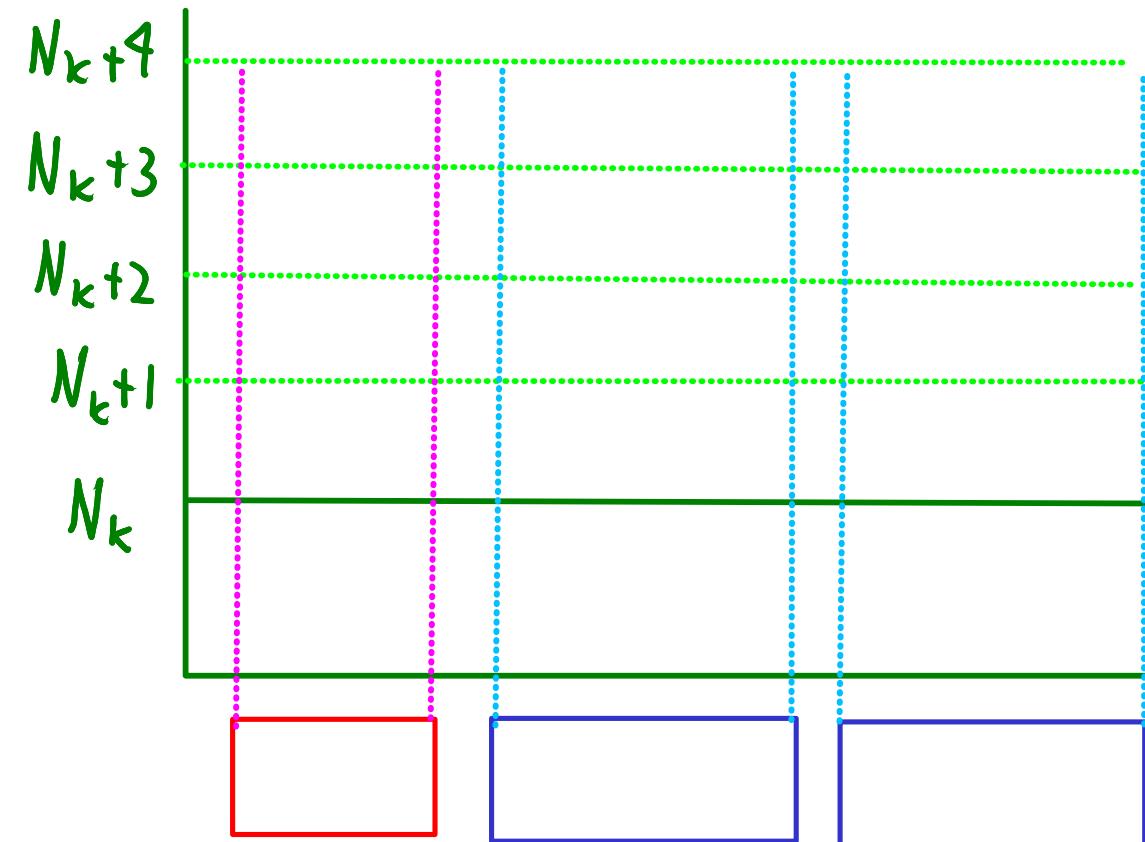


$F_0 \setminus RC(F)$

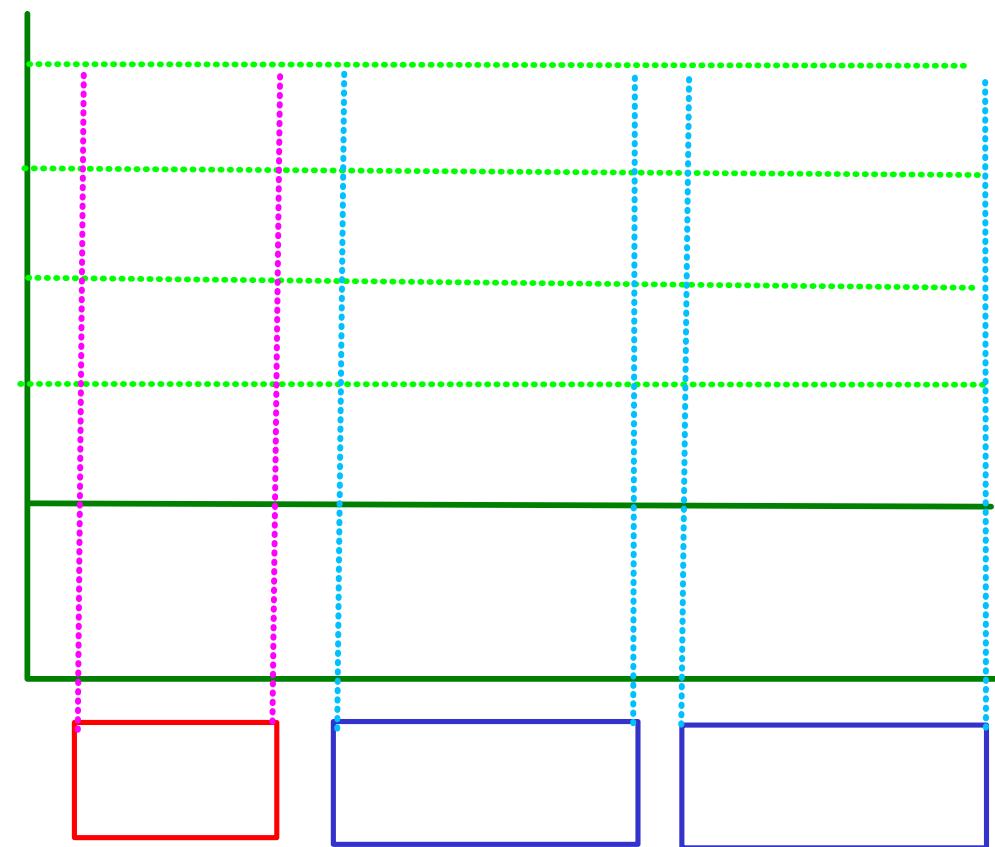


$F_1 \setminus RC(F)$

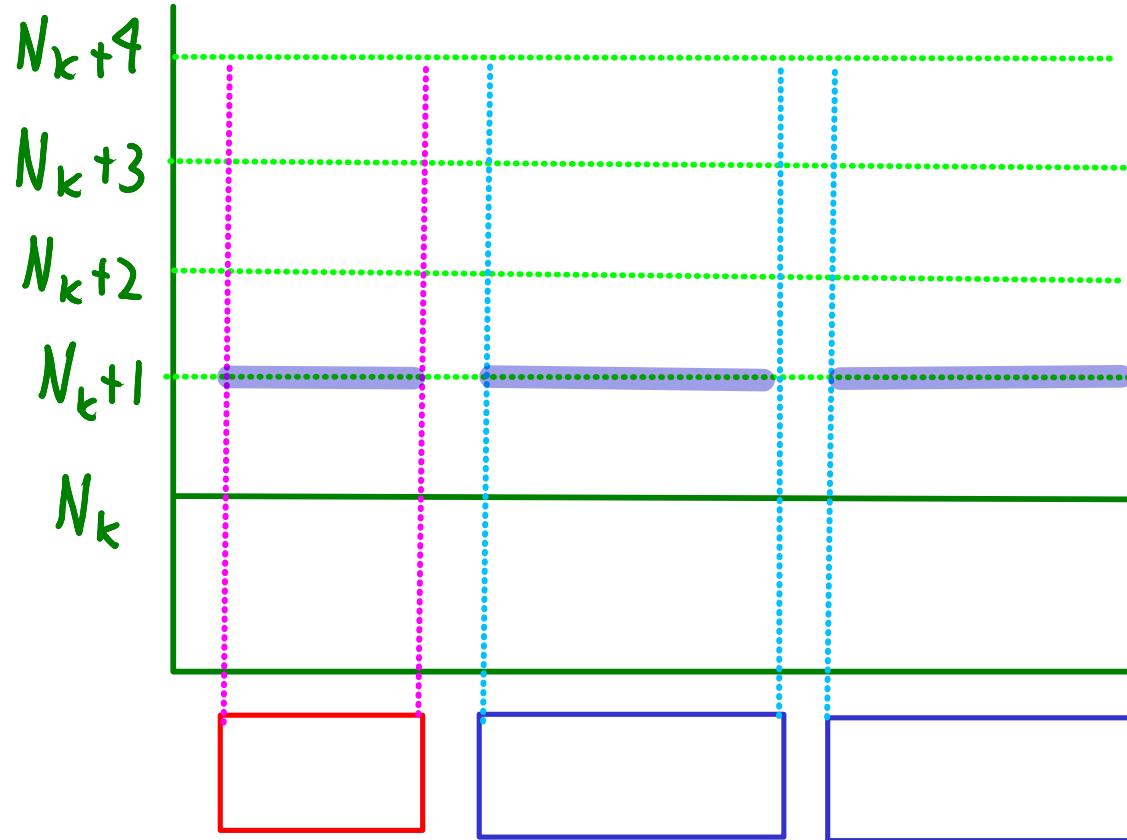
$A^F$



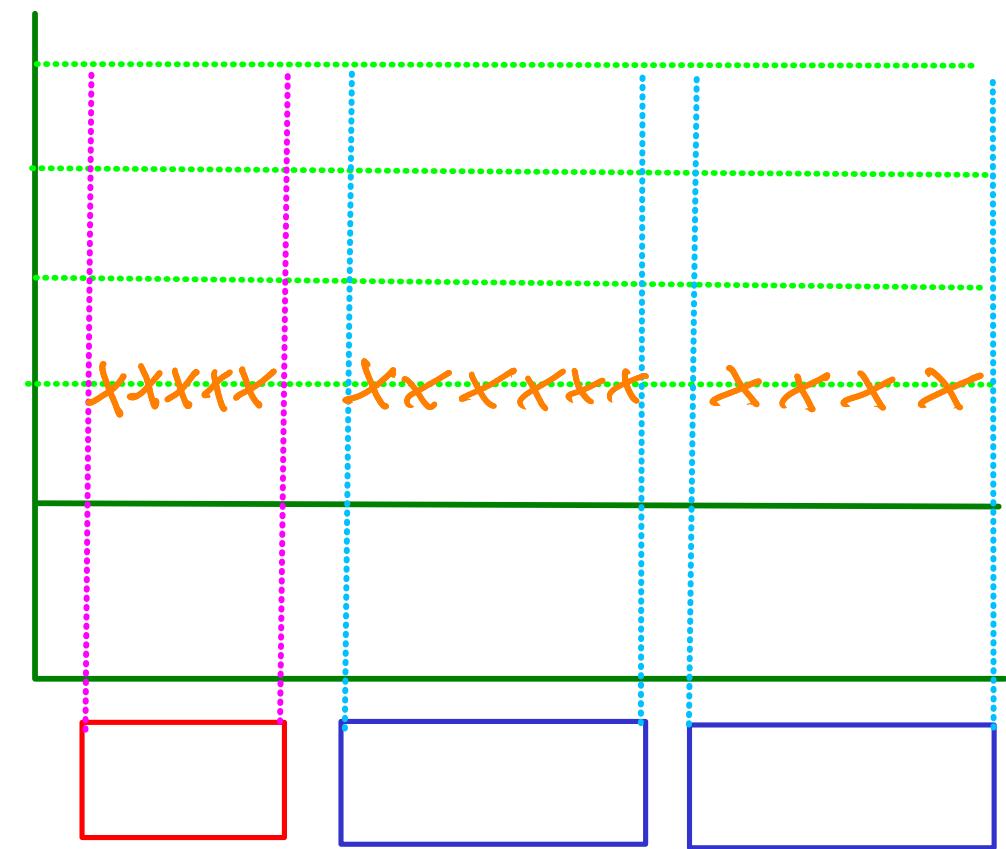
$B^F$



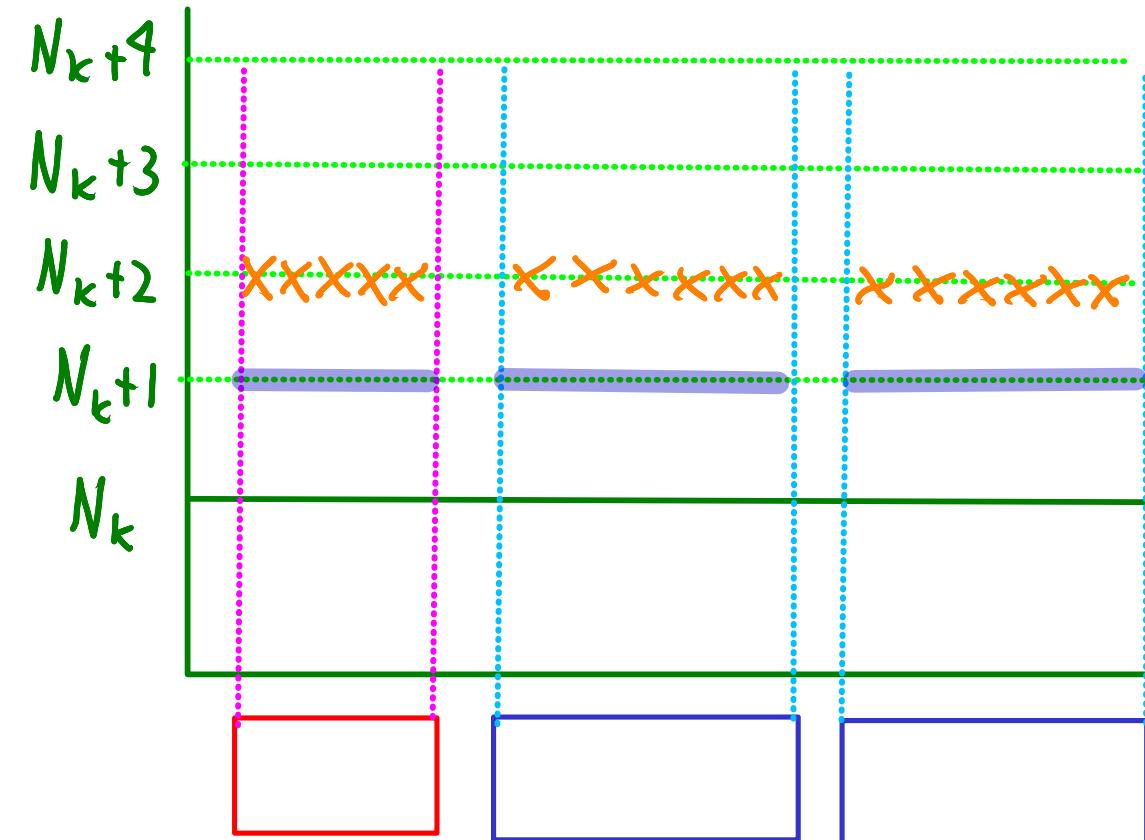
$A^F$



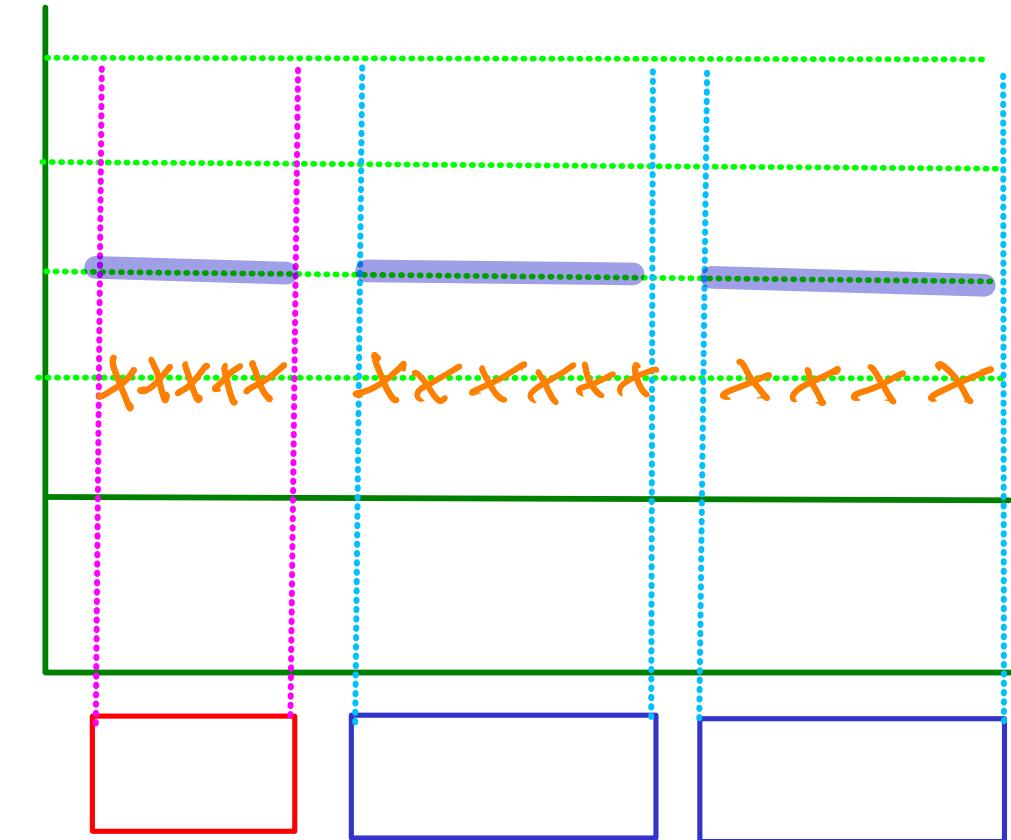
$B^F$



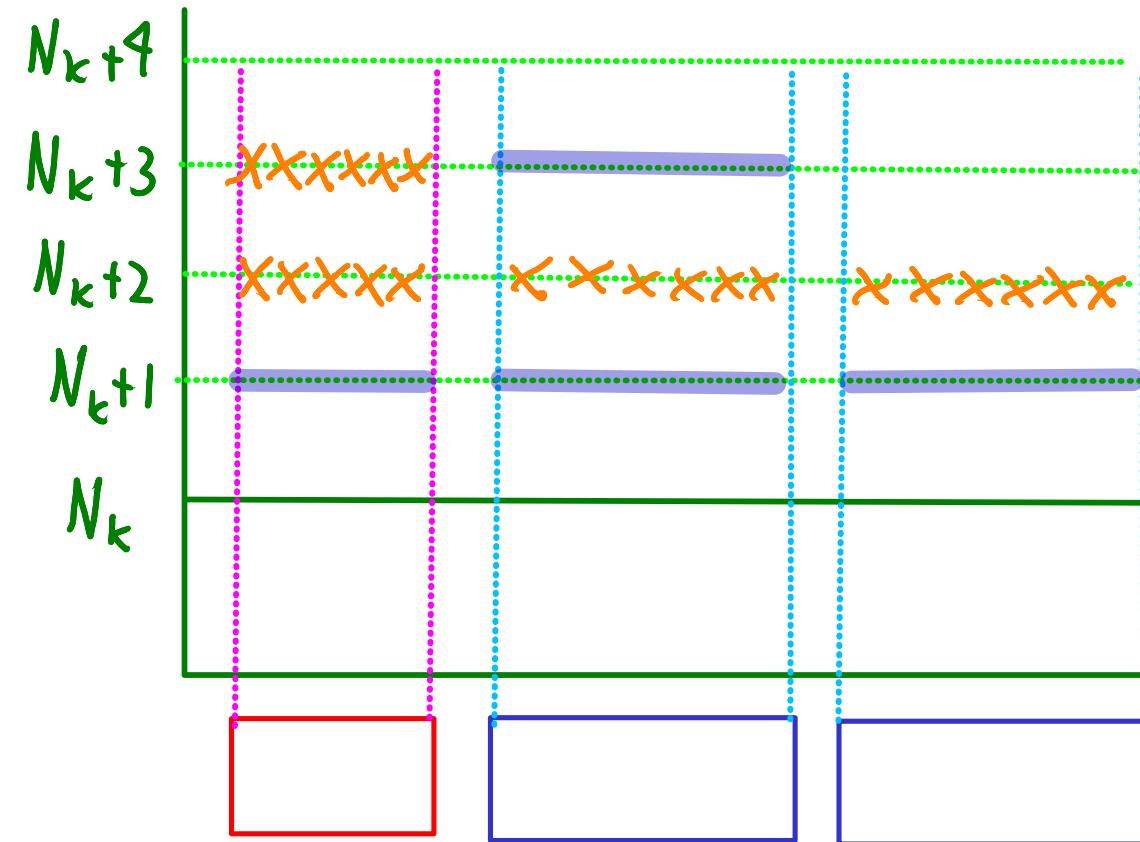
$A^F$



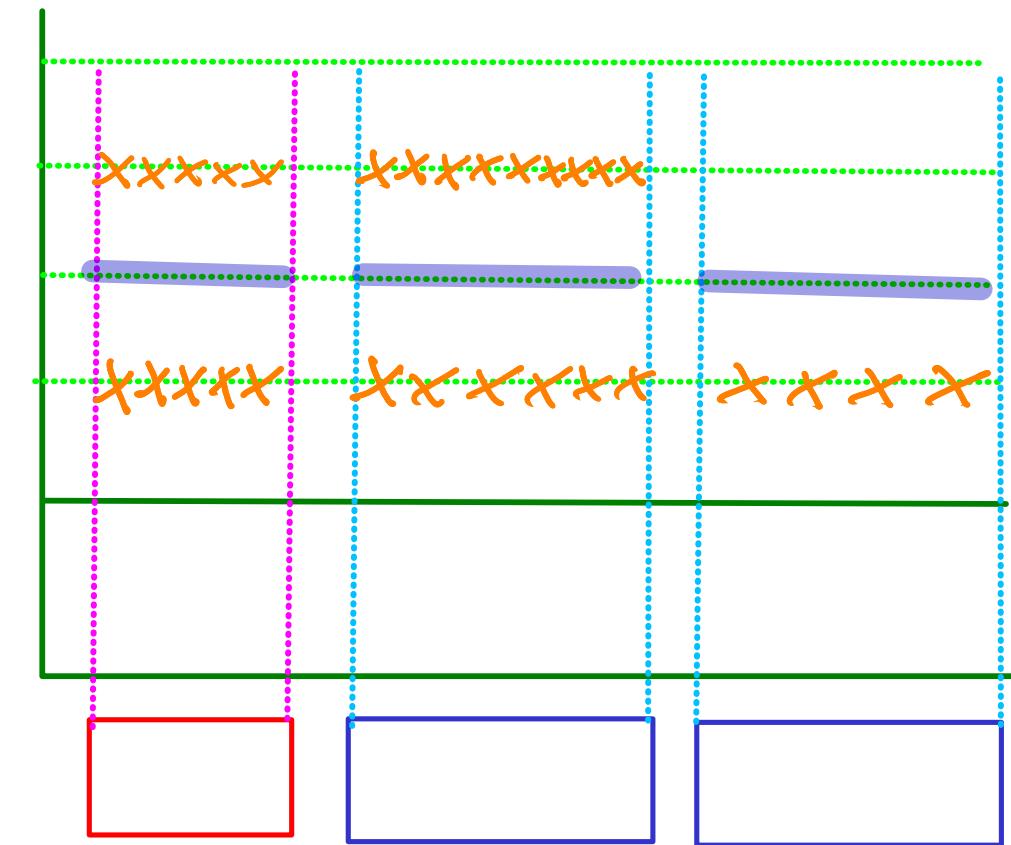
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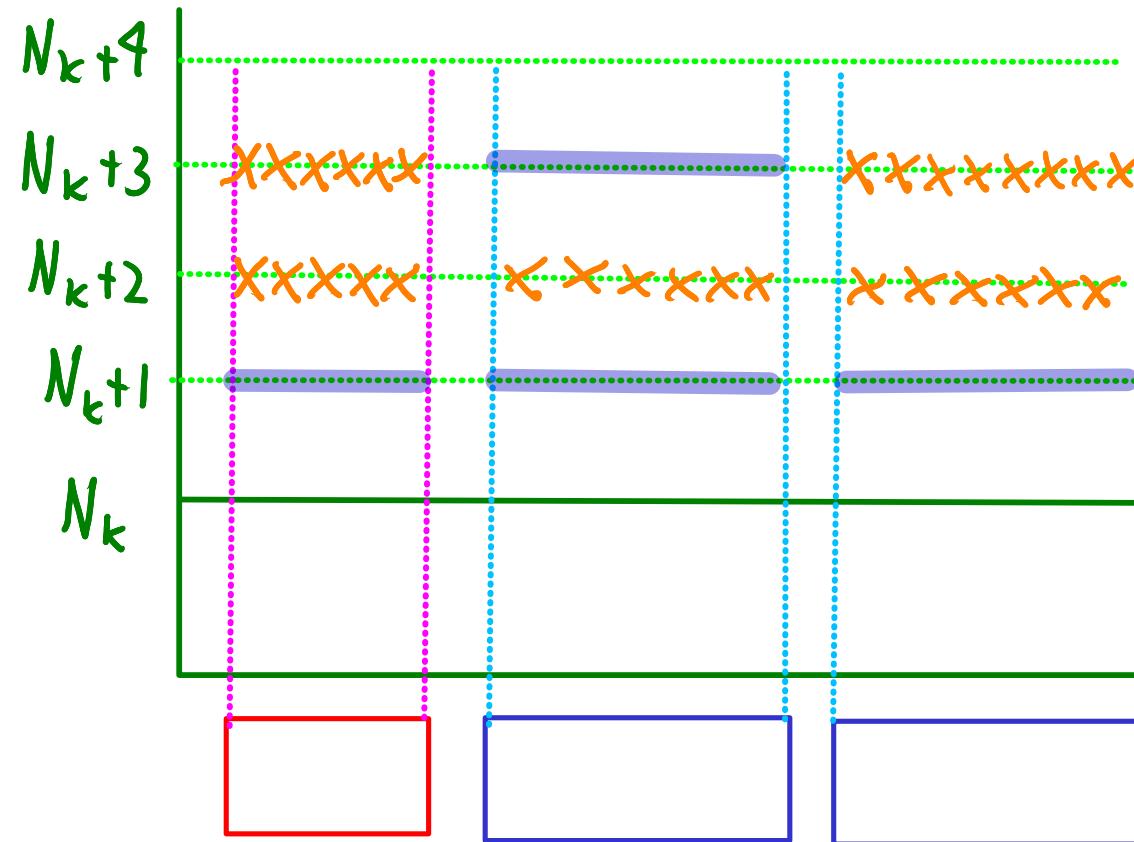
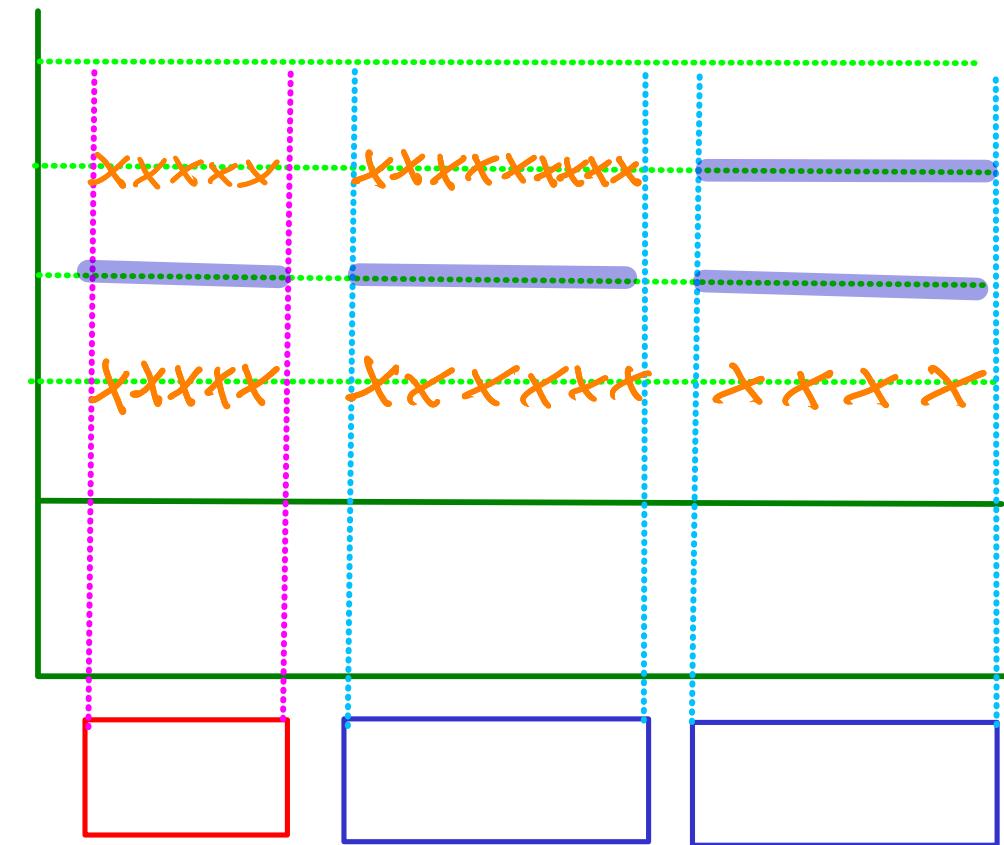


$A^F$

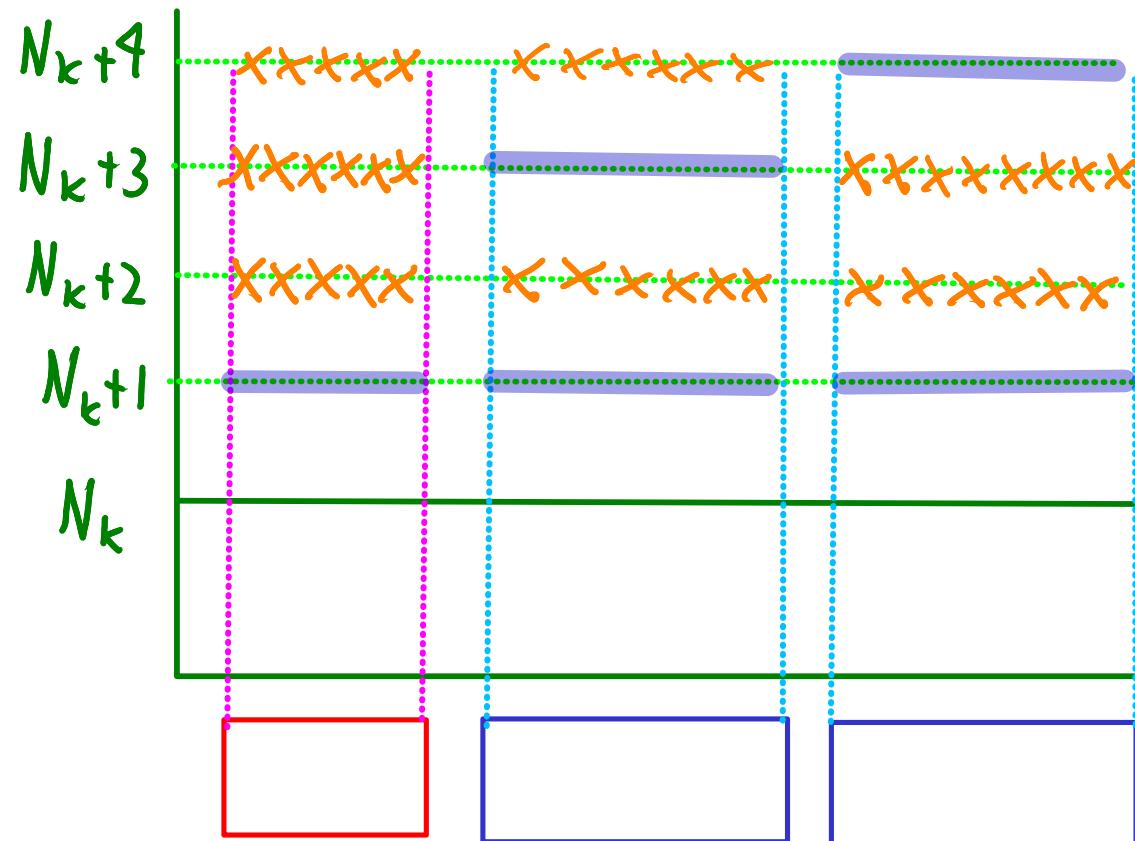


$B^F$

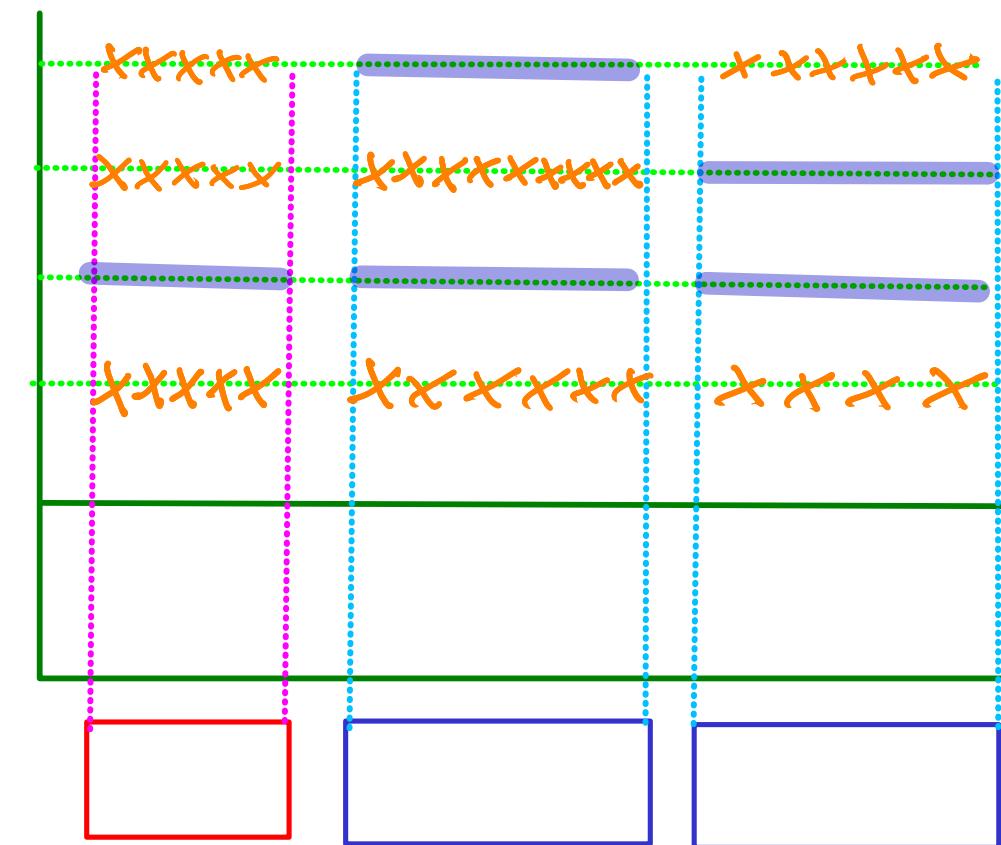


$A^F$  $B^F$ 

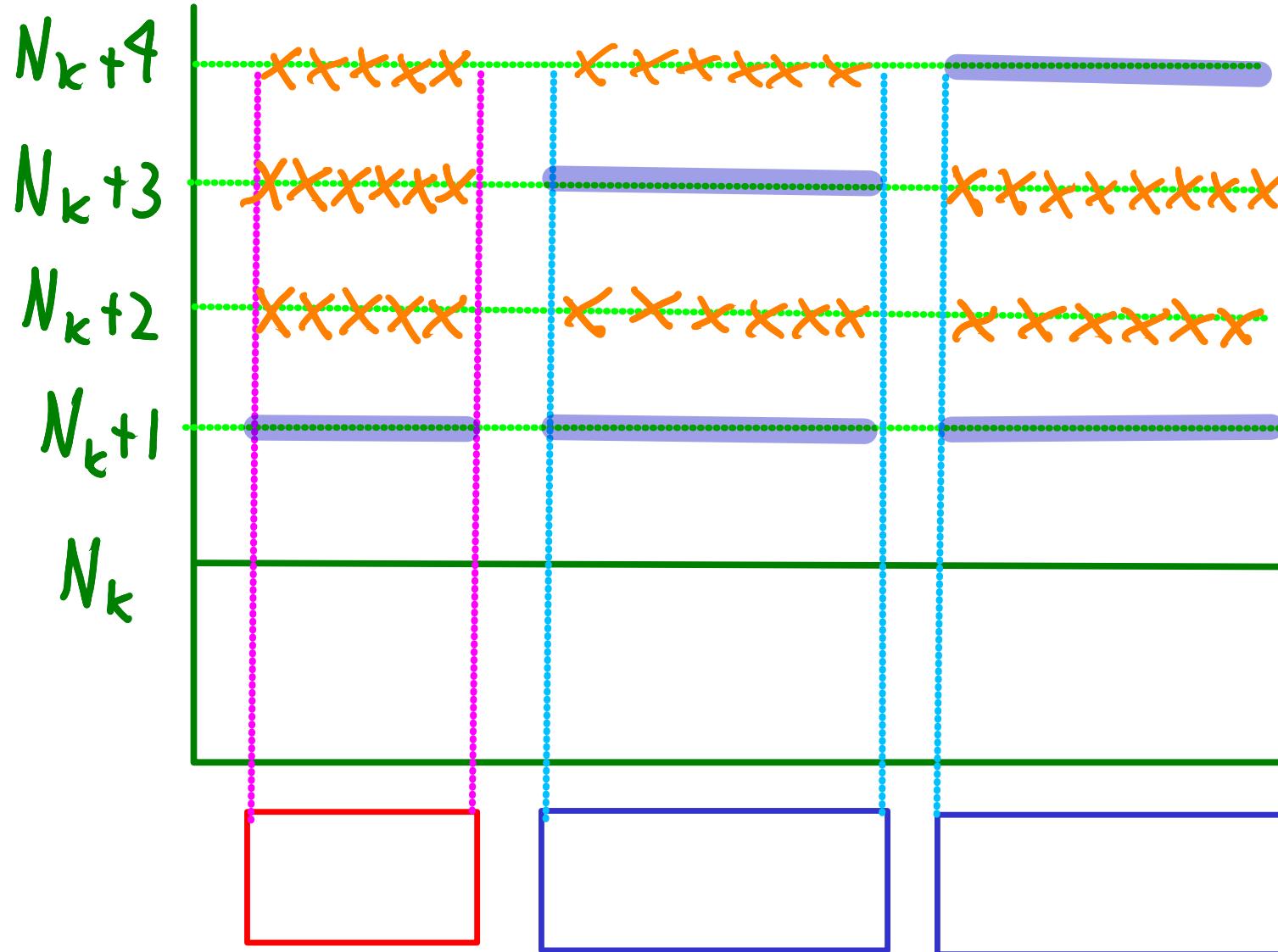
$A^F$



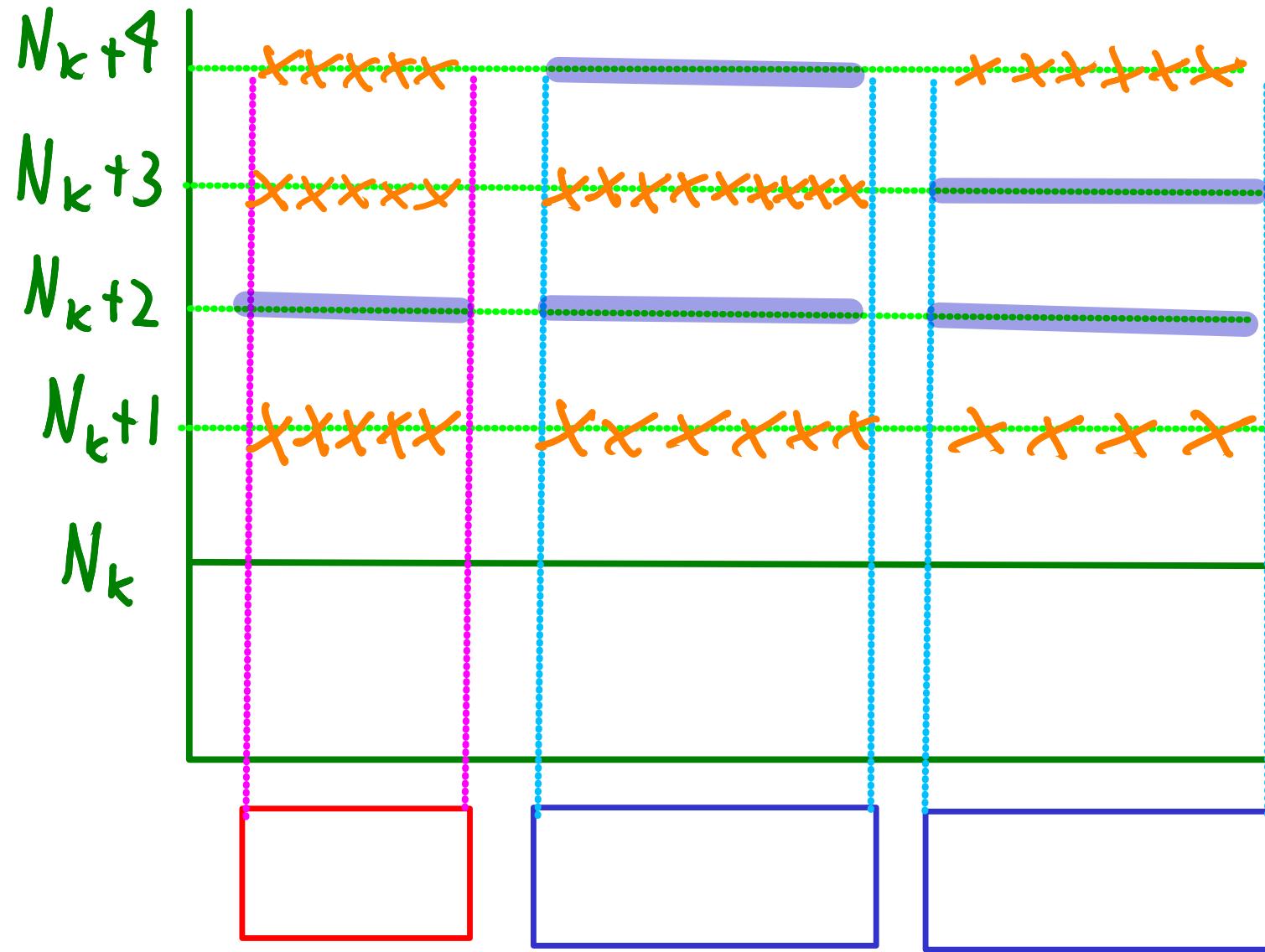
$B^F$



$A^F$



$\beta^F$

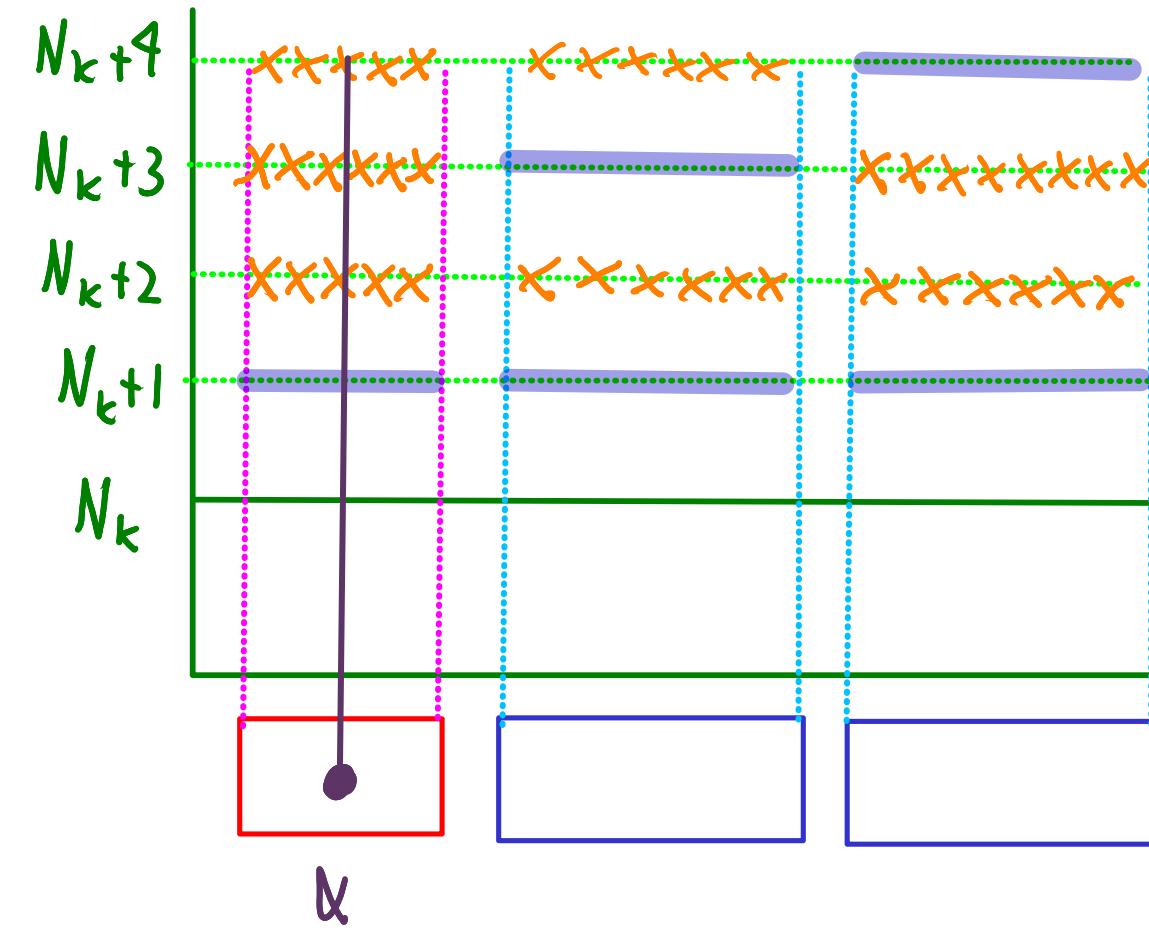
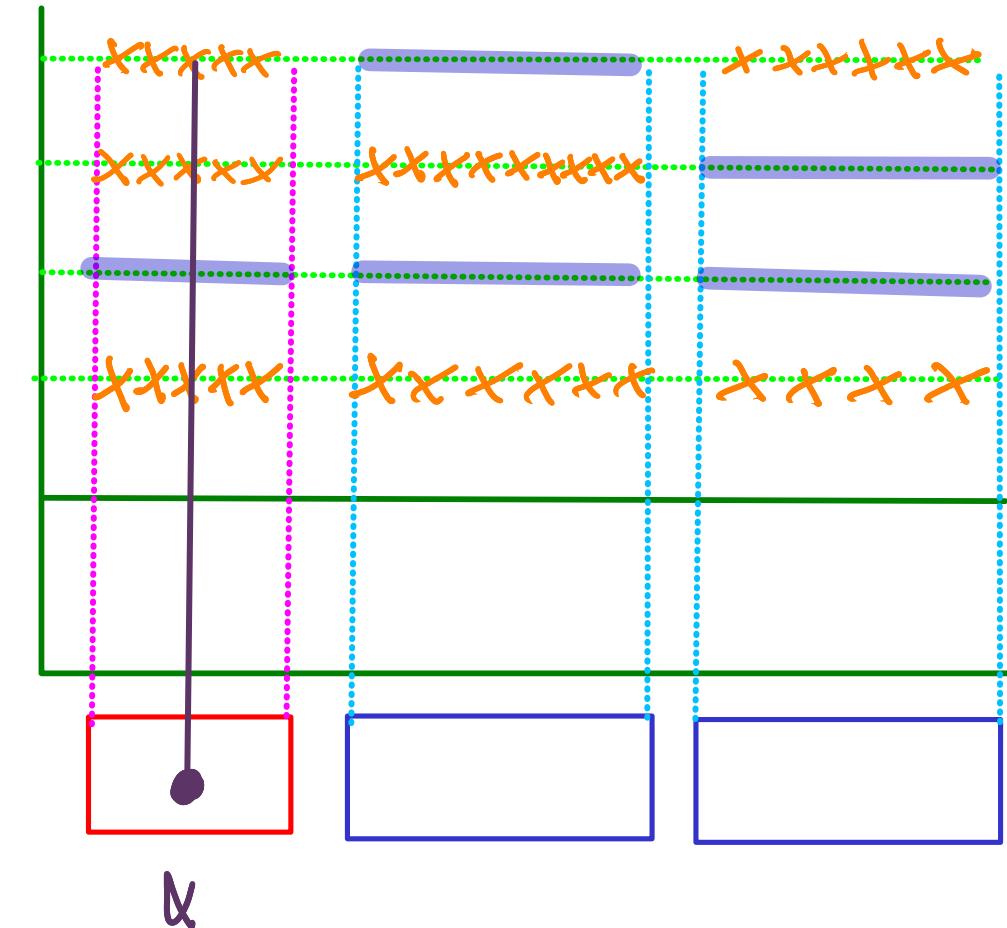


In words:

1) If  $\alpha \in RCF$ )

$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}\}$$

$A^F$  $B^F$ 

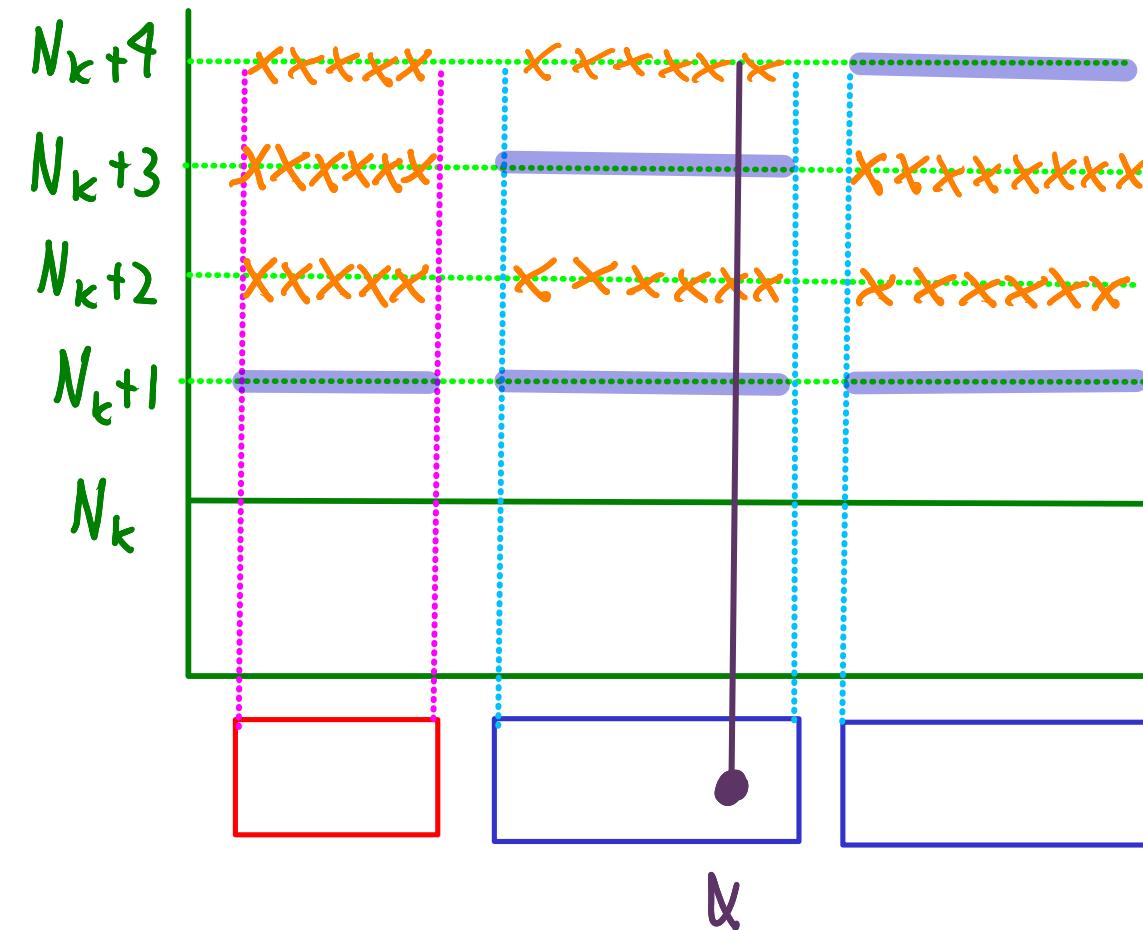
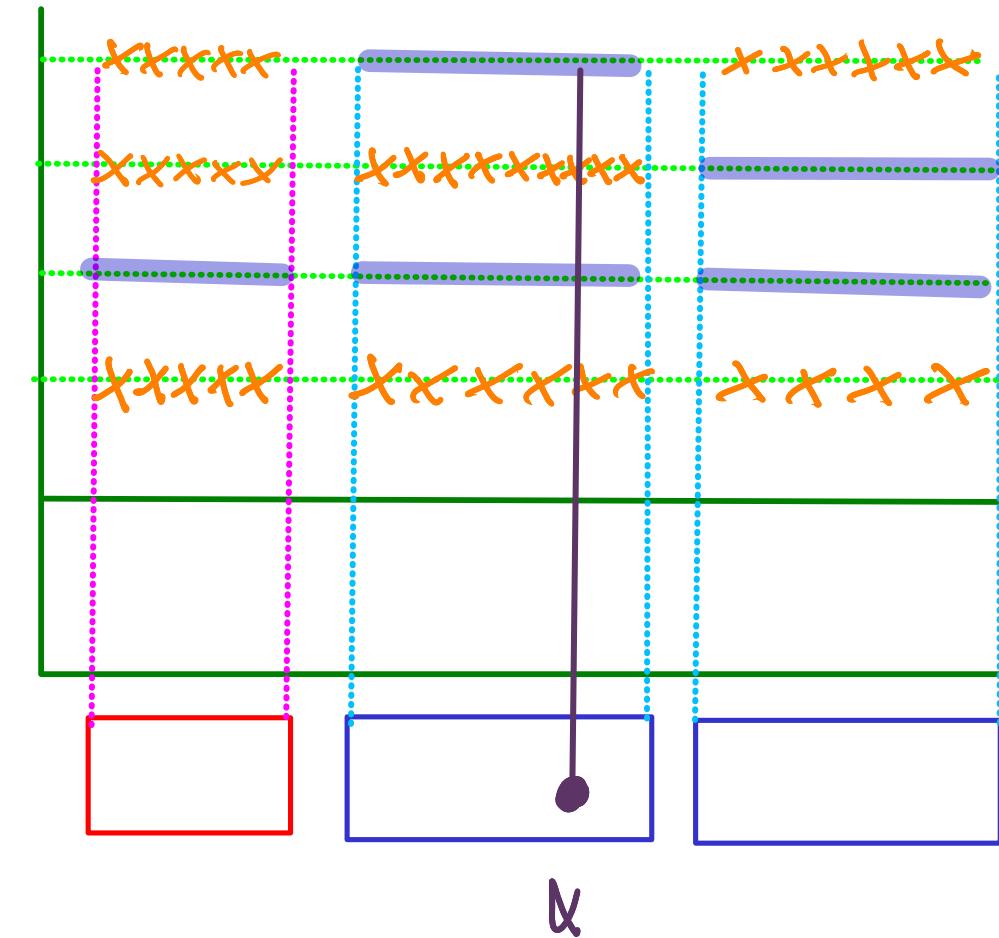
$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}\}$$

2) If  $\alpha \in F_0 \setminus R(F)$

$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}, N_{k+3}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}, N_{k+4}\}$$

$A^F$  $B^F$ 

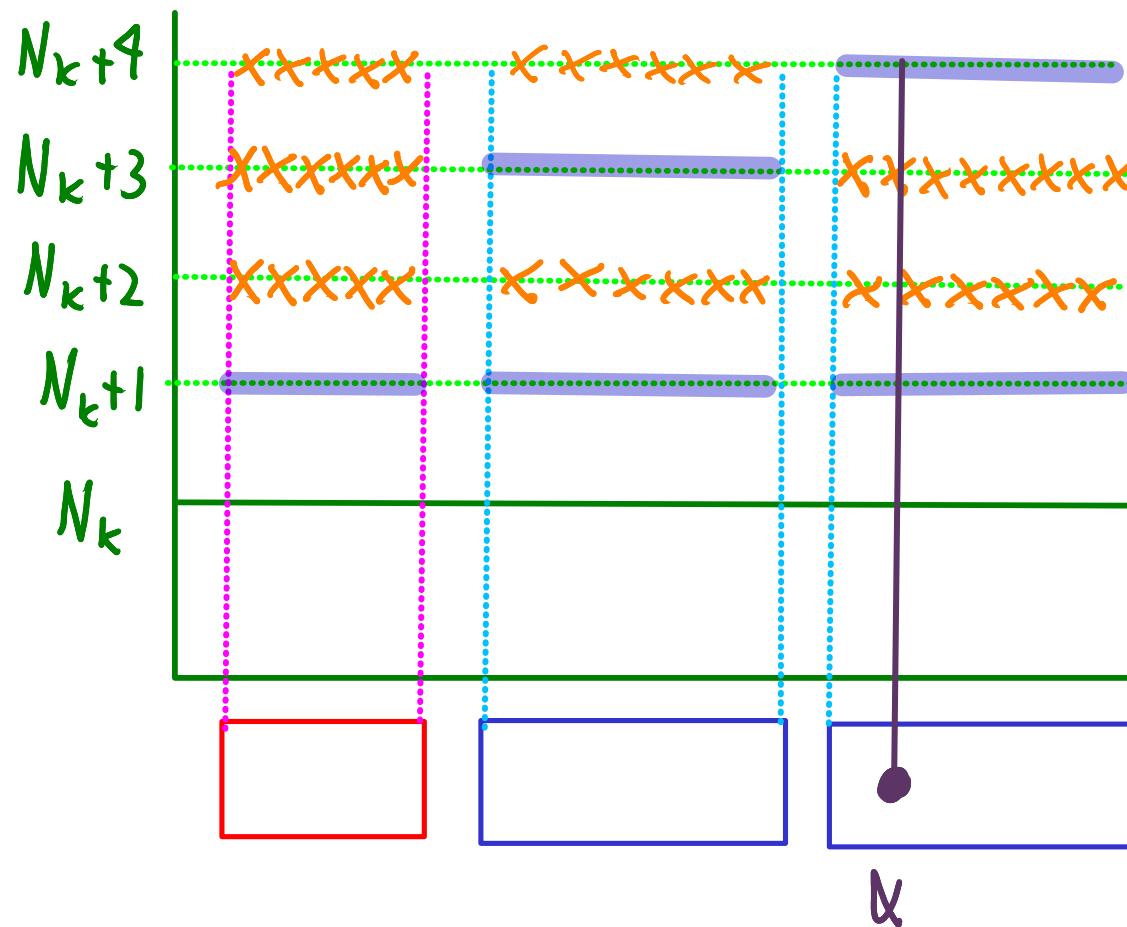
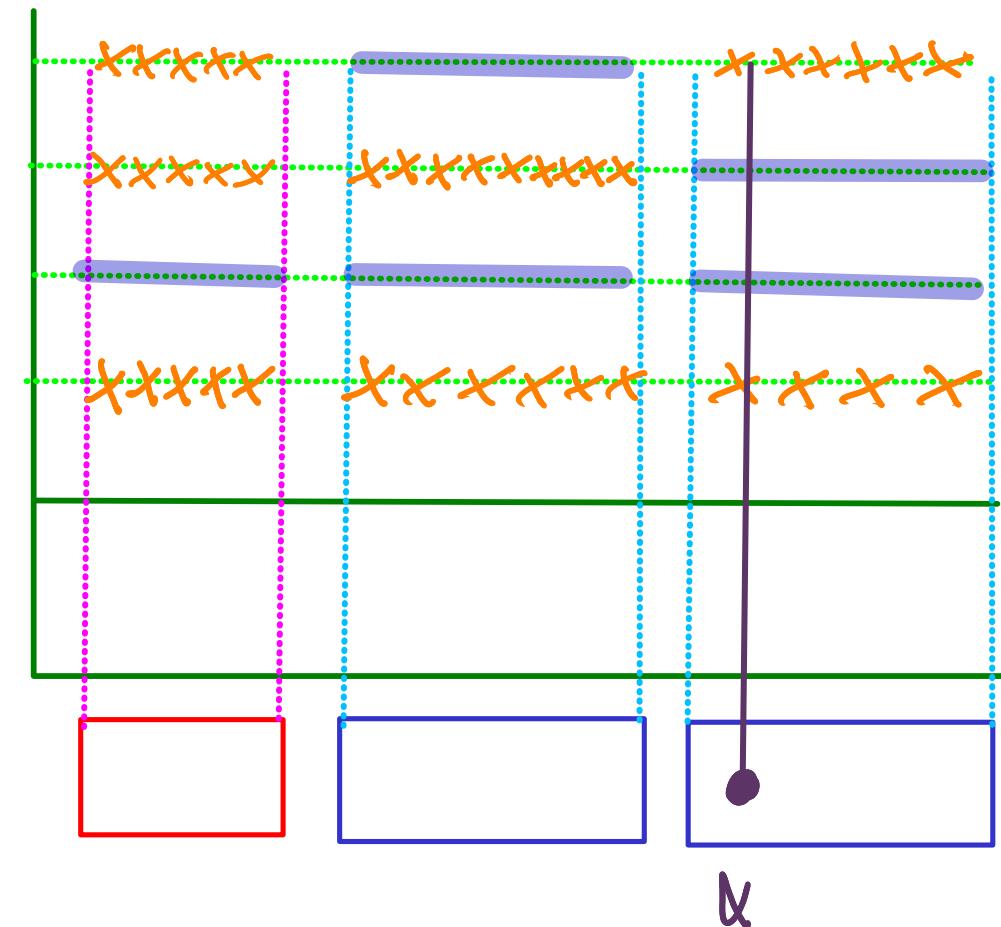
$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}, N_{k+3}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}, N_{k+4}\}$$

3) If  $\alpha \in F_i \setminus R(F)$

$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}, N_{k+4}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}, N_{k+3}\}$$

$A^F$  $B^F$ 

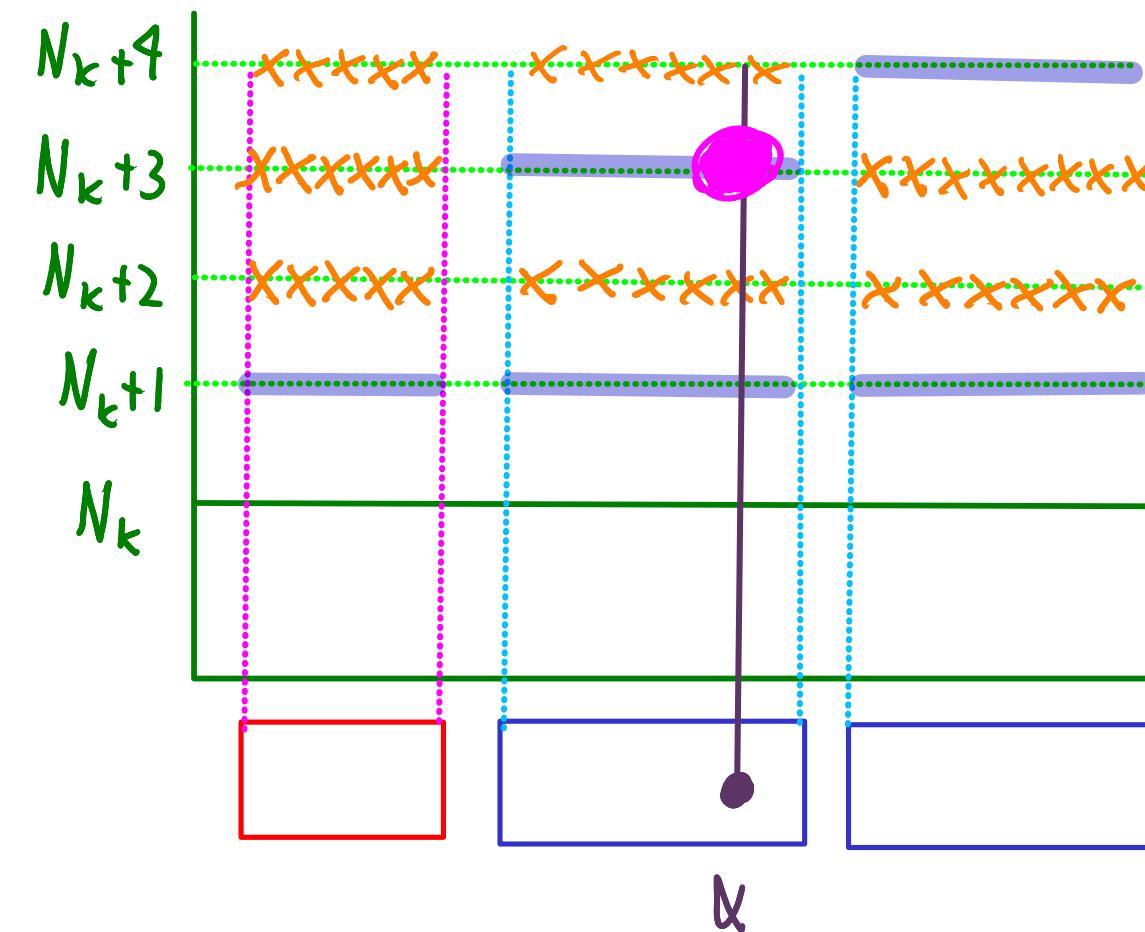
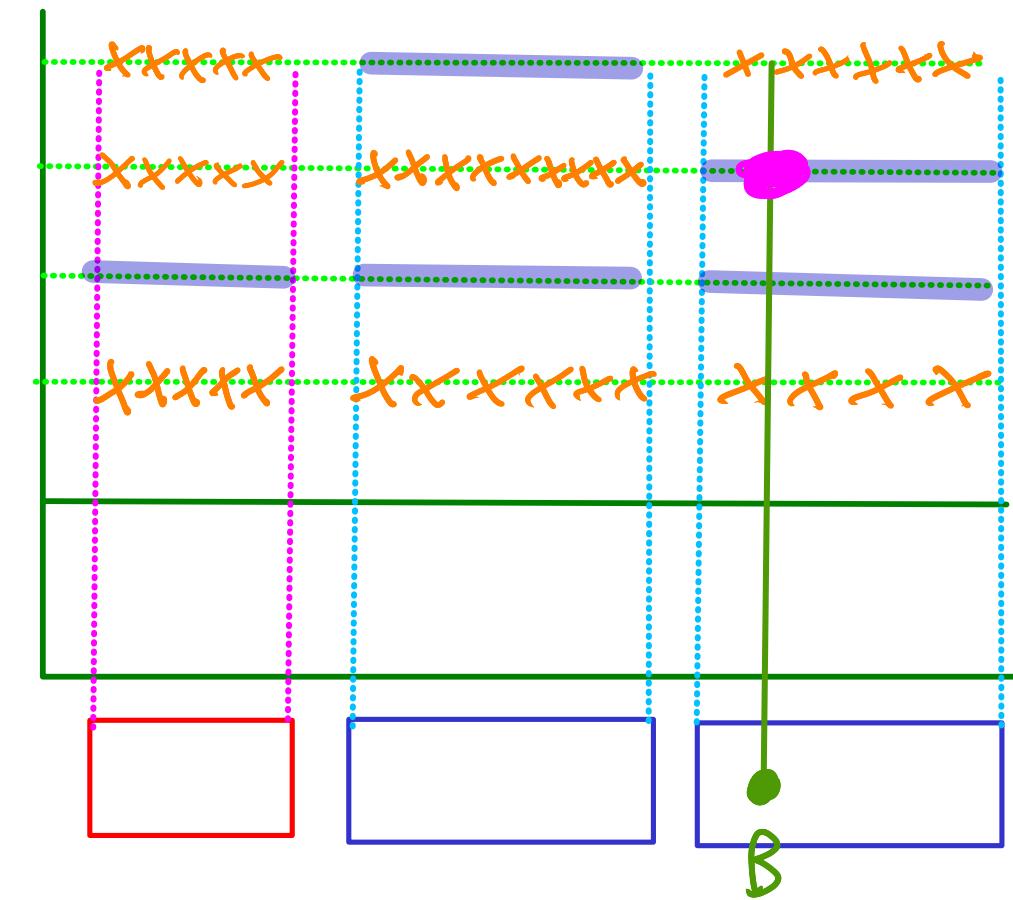
$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}, N_{k+4}\}$$

$$B_\alpha^F = B_\alpha^{F_0} \cup \{N_{k+2}, N_{k+3}\}$$

The point is that if :

$$\alpha \in F_0 \setminus RCF \quad \text{and} \quad \beta \in F_1 \setminus RCF$$

Then:

$A^F$  $B^F$ 

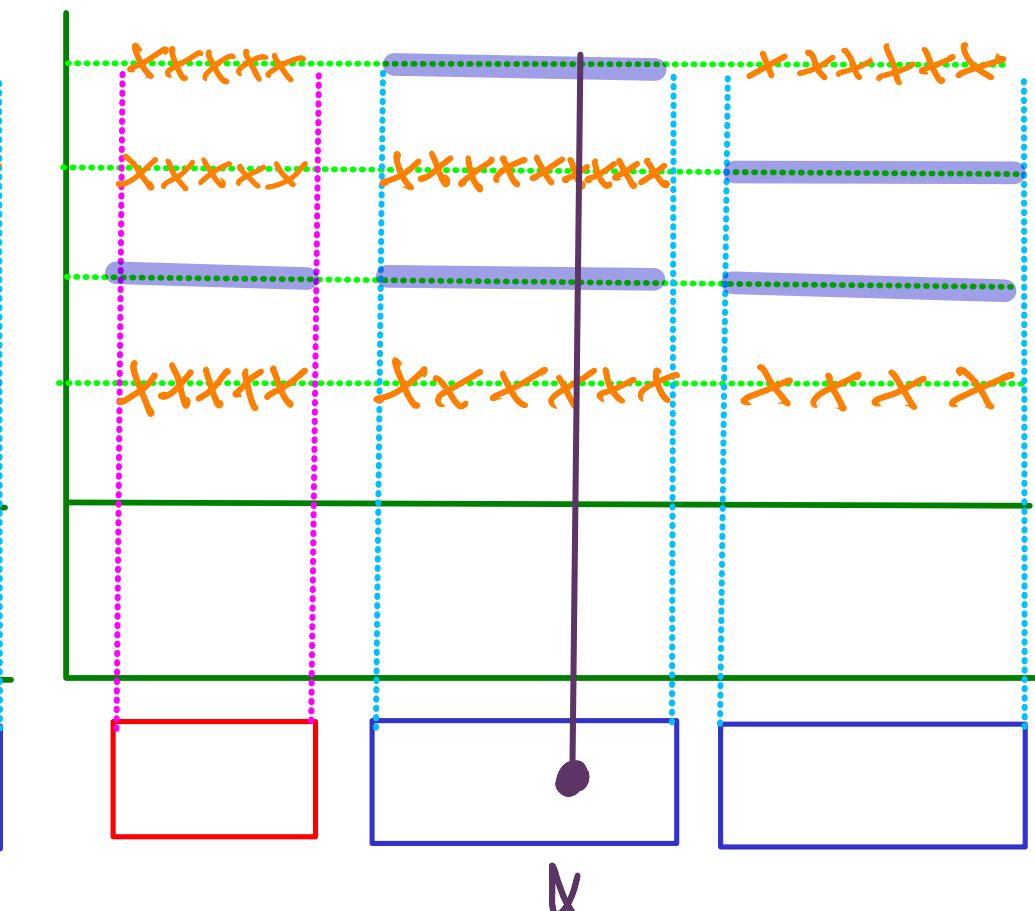
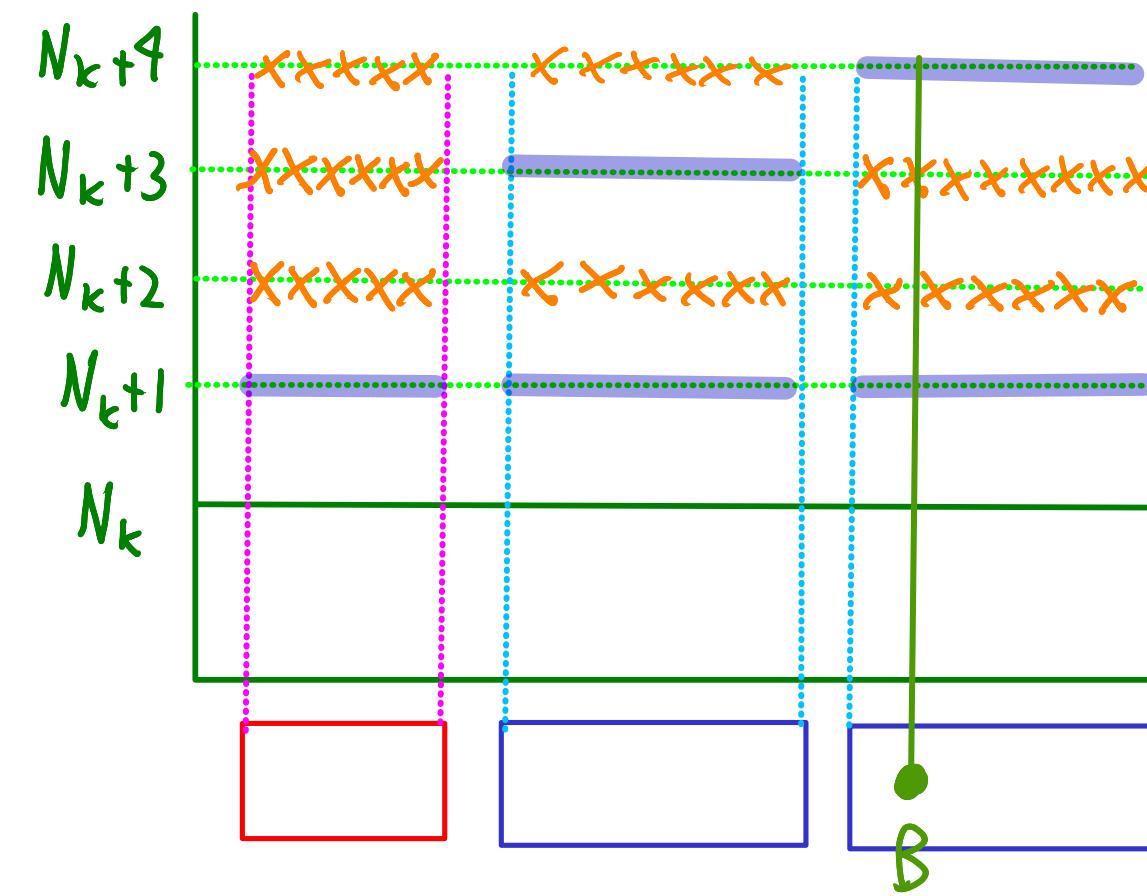
$$A_\alpha^F = A_\alpha^{F_0} \cup \{N_{k+1}, N_{k+3}\}$$

$$B_B^F = B_B^{F_0} \cup \{N_{k+2}, N_{k+3}\}$$

$N_{k+3} \in A_\alpha^F \cap B_B^F$

$A^F$

$B^F$



$$A_\beta^F = A_\beta^{F_1} \cup \{N_{k+1}, N_{k+4}\}$$

$$B_\alpha^F = B_\alpha^{F_1} \cup \{N_{k+2}, N_{k+4}\}$$

$N_{k+4} \in A_\beta^F \cap B_\alpha^F$

This finishes the recursion  
construction

Now, for every  $\alpha < \omega_1$ , define:

$$A_\alpha = \bigcup \{ A_\alpha^F \mid \alpha \in F \in \mathcal{F} \}$$

$$B_\alpha = \bigcup \{ B_\alpha^F \mid \alpha \in F \in \mathcal{F} \}$$

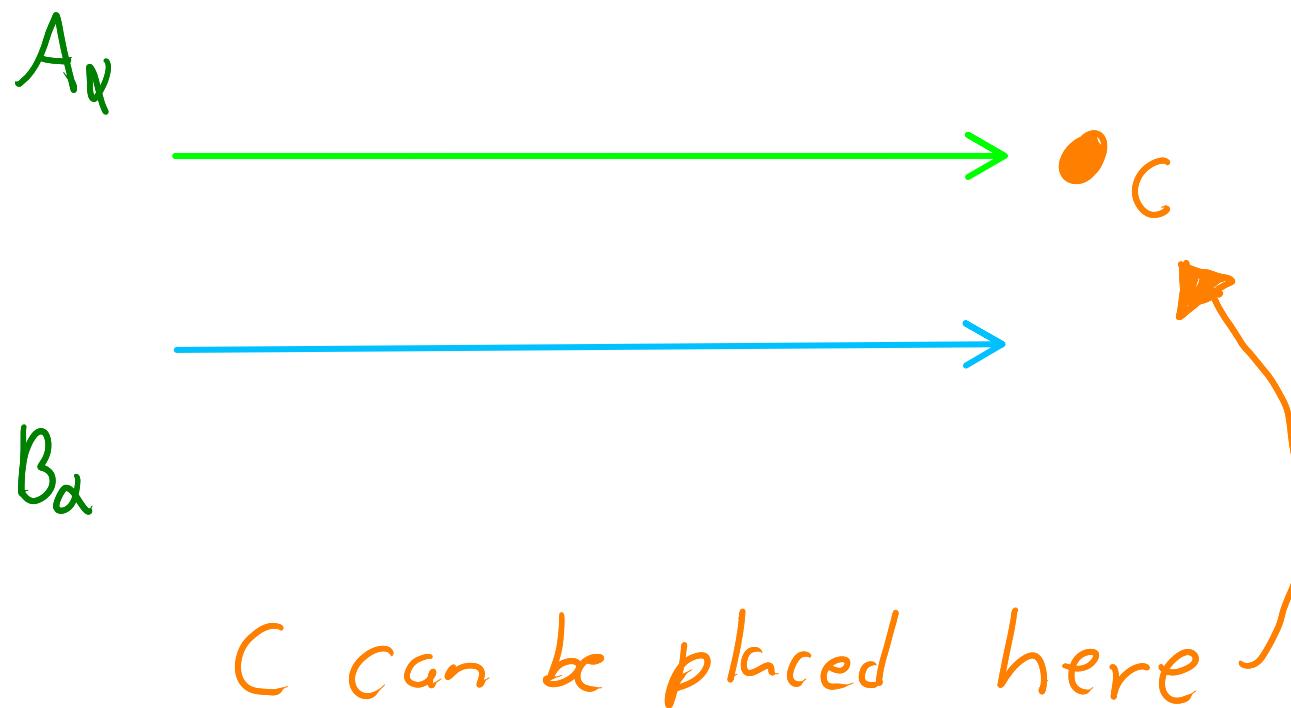
## Exercise

Prove that  $(\langle A_\alpha \rangle, \langle B_\alpha \rangle)$  is a pregap

We now prove it is a gap. Assume  
there is  $C \subseteq \omega$  such that for every  
 $\alpha < \omega_1$ :

- 1)  $A_\alpha \subseteq^* C$
- 2)  $B_\alpha \cap C \text{ is finite}$

We now prove it is a gap. Assume there is  $C \subseteq w$  such that:



For every  $\alpha < \omega_1$ , we can find  $l_\alpha \in w$   
such that:

$$A_\alpha \setminus l_\alpha \subseteq C$$

$$B_\alpha \setminus l_\alpha \subseteq w \setminus C$$

By a counting argument, we can  
 $\lambda < \beta$  and  $\lambda < \omega$  such that:

$$1) \lambda = \lambda_\alpha = \lambda_\beta$$

$$2) A_\alpha \cap l = A_\beta \cap l$$

$$B_\alpha \cap l = B_\beta \cap l$$

Now, let  $k \in \omega$  be the first such  
that there is  $F \in \mathcal{F}$  such that

$$\alpha, \beta \in F$$

Now, let  $kew$  be the first such  
that there is  $F \in \mathcal{F}$  such that

$$\alpha, \beta \in F$$

(In other words,  $k$  is the "distance"  
between  $\alpha$  and  $\beta$ , as defined in  
the previous lecture)

Claim

$\alpha \in F_0 \setminus R(F)$  and  $\beta \in F_1 \setminus R(F)$

Claim

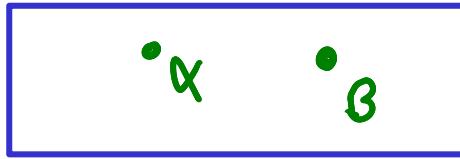
$\alpha \in F_0 \setminus R(F)$  and  $\beta \in F_1 \setminus R(F)$

If  $\beta \in F_0$  then:



$\alpha$  will also be in  $F_0$  ( $\alpha < \beta$ ), but this contradicts the minimality of  $k$

If  $\alpha \in F_i$  then:



$\beta$  will also be in  $F_i$  ( $\alpha < \beta$ ), but this contradicts the minimality of  $k$

This finishes the proof of the claim

By the construction, we get that

there is  $i \in A_\alpha \cap B_\beta$ .

claim

$$l \leq i$$

claim

$$l \leq i$$

Assume  $i < l$ . Now:

$$i \in A_\alpha \cap l = A_\beta \cap l$$

$$\Rightarrow i \in A_\beta \cap B_\beta$$

which is a contradiction

Recall:

$$i \in A_\alpha \cap B_\beta$$

$$A_\alpha \cap l = A_\alpha \cap l$$

Finally, recall that:

$$A_\alpha \setminus l \subseteq C$$

$$B_\beta \setminus l \subseteq w \setminus C$$

so  $i \in A_\alpha \setminus l \Rightarrow i \in C$

$$i \in B_\beta \setminus l \Rightarrow i \in w \setminus C$$

which is obviously a contradiction



# Constructing an Aronszajn tree

Def

An Aronszajn tree is a tree such that:

- 1) It has height  $\omega_1$ ,
- 2) Its levels are countable
- 3) It has no cofinal branches